# **6.033 Spring 2015**Lecture #7

- Monolithic kernels vs. Microkernels
- Virtual Machines

### Enforcing Modularity via Virtualization

in order to enforce modularity + build an effective operating system

 programs shouldn't be able to refer to (and corrupt) each others' memory

virtual memory

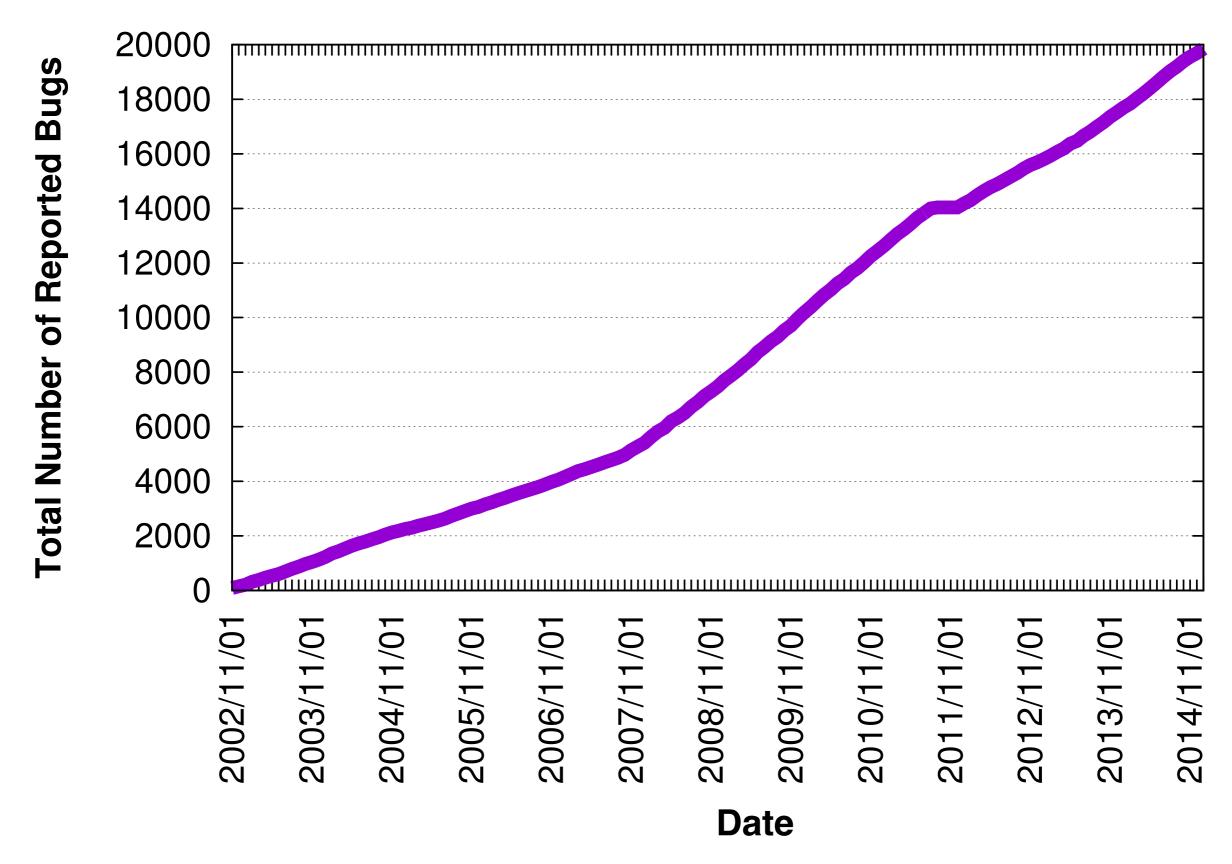
programs should be able to communicate

bounded buffers
(virtualize communication links)

3. programs should be able to **share a CPU** without one program halting the progress of the others

threads (virtualize processors)

today: can we rely on the kernel to work properly?



source: bugzilla.kernel.org, count of all bugs currently marked NEW, ASSIGNED, REOPENED, RESOLVED, VERIFIED, or CLOSED, by creation date

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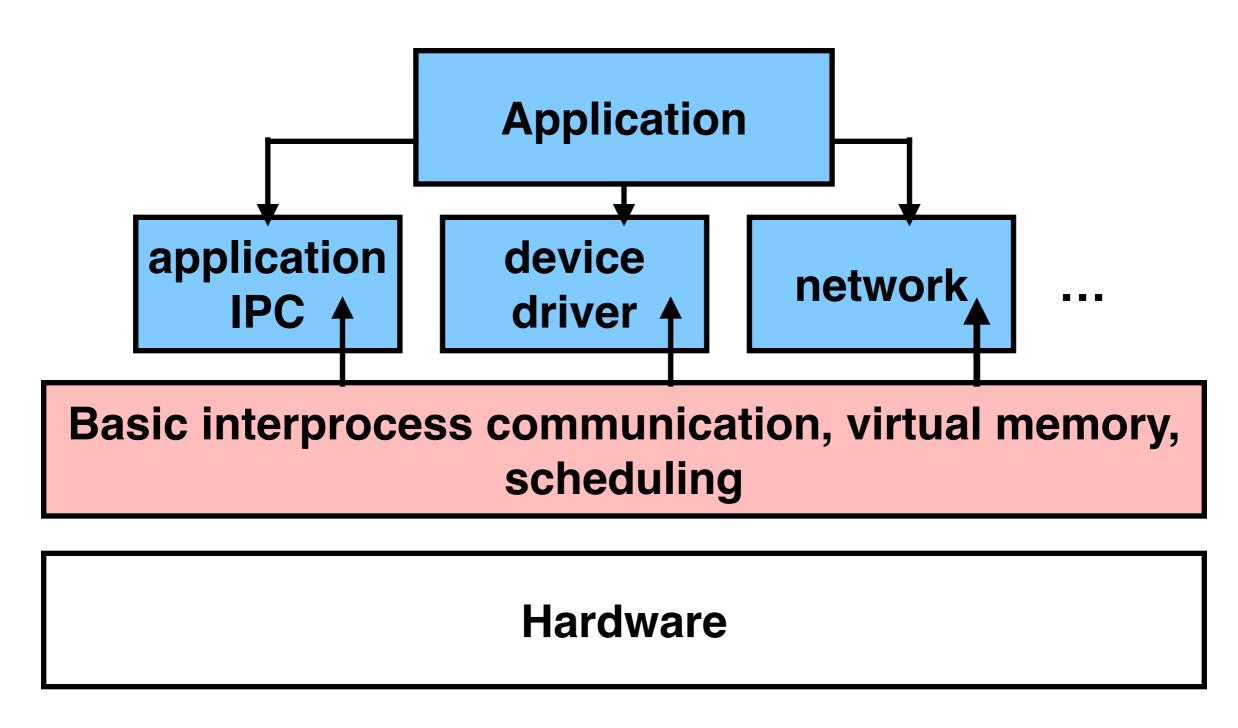
# monolithic kernels: no enforced modularity within the kernel itself

**Application** 

Basic interprocess communication, virtual memory, scheduling, file server, device drivers, network, ...

#### **Hardware**

microkernels: enforce modularity by putting subsystems in user programs



problem: how do we deal with bugs in the Linux kernel without redesigning Linux from scratch?

#### Virtual Machines

one approach: VMM runs as user program inside host OS

#### host OS

virtual machine running guest OS

virtual machine running guest OS

virtual machine monitor (VMM)

#### Virtual Machines

another approach: VMM runs in kernel-mode on hardware

virtual machine running guest OS

virtual machine running guest OS

virtual machine monitor (VMM)

physical hardware

#### guest OS

#### guest OS

#### virtual hardware

U/K PTR page table

#### virtual hardware

U/K PTR page table

#### virtual machine monitor (VMM)

#### physical hardware

U/K, PTR, page table, ...

#### VMM's goal: virtualize hardware

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#### guest OS

#### guest OS

## guest virtual address

#### virtual hardware

U/K
PTR
page table

#### virtual hardware

U/K
PTR
page table

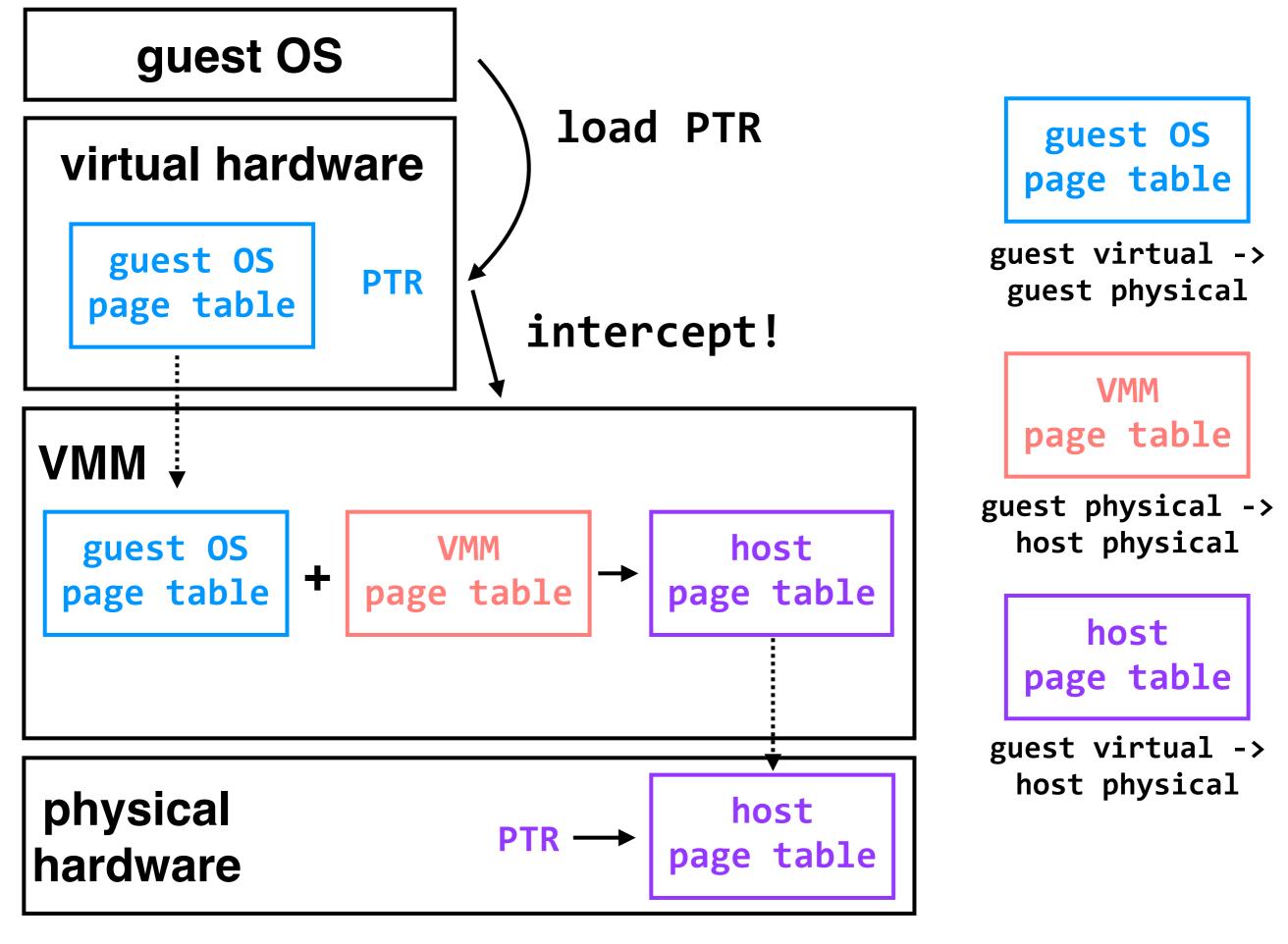
guest physical
 address

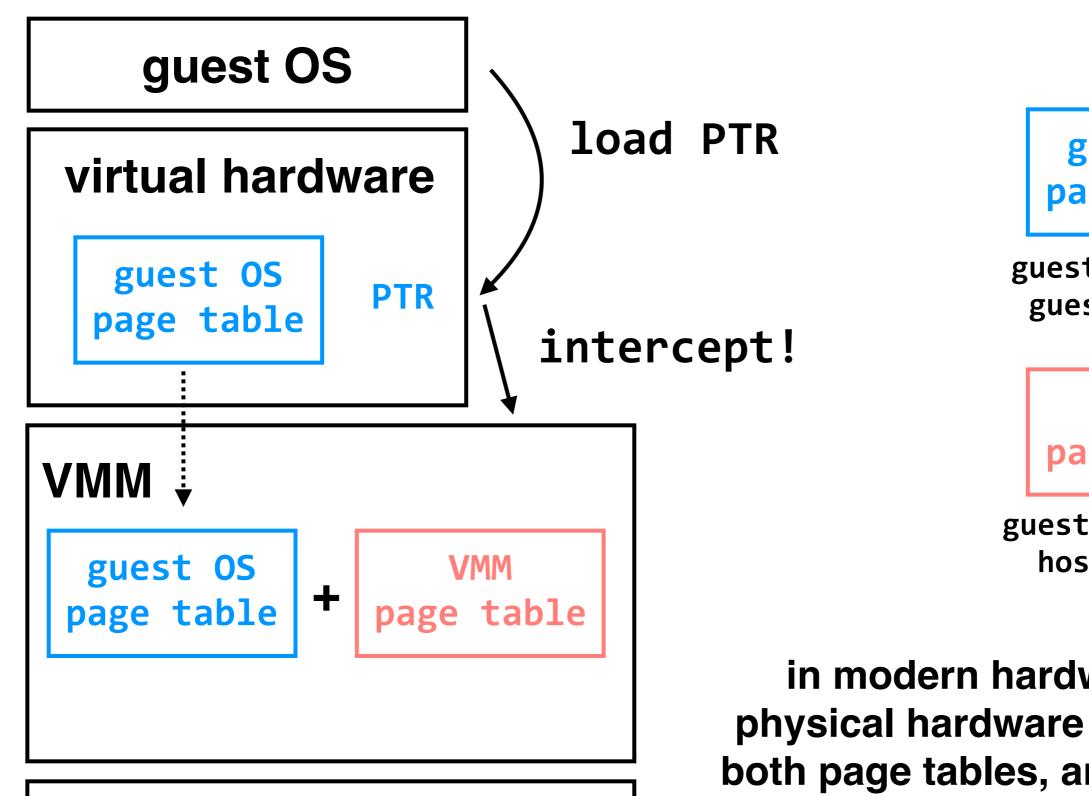
#### virtual machine monitor (VMM)

#### physical hardware

U/K, PTR, page table, ...

host physical address





**PTR** 

physical

hardware

guest OS
page table

guest virtual ->
 guest physical

VMM page table

guest physical ->
 host physical

in modern hardware, the physical hardware is aware of both page tables, and performs the translation from guest virtual to host physical itself

#### Kernel Structure

**Monolithic kernels** provide no enforced modularity within the kernel. **Microkernels** do, but redesigning monolithic kernels as microkernels is challenging.

#### Virtual Machines

Virtual machines allow us to run multiple **isolated**OSes on a single physical machine, similar to how we used an OS to run multiple programs on a single CPU.
VMs must handle the challenges of virtualizing the hardware (examples: virtualizing memory, the U/K bit, and disk).