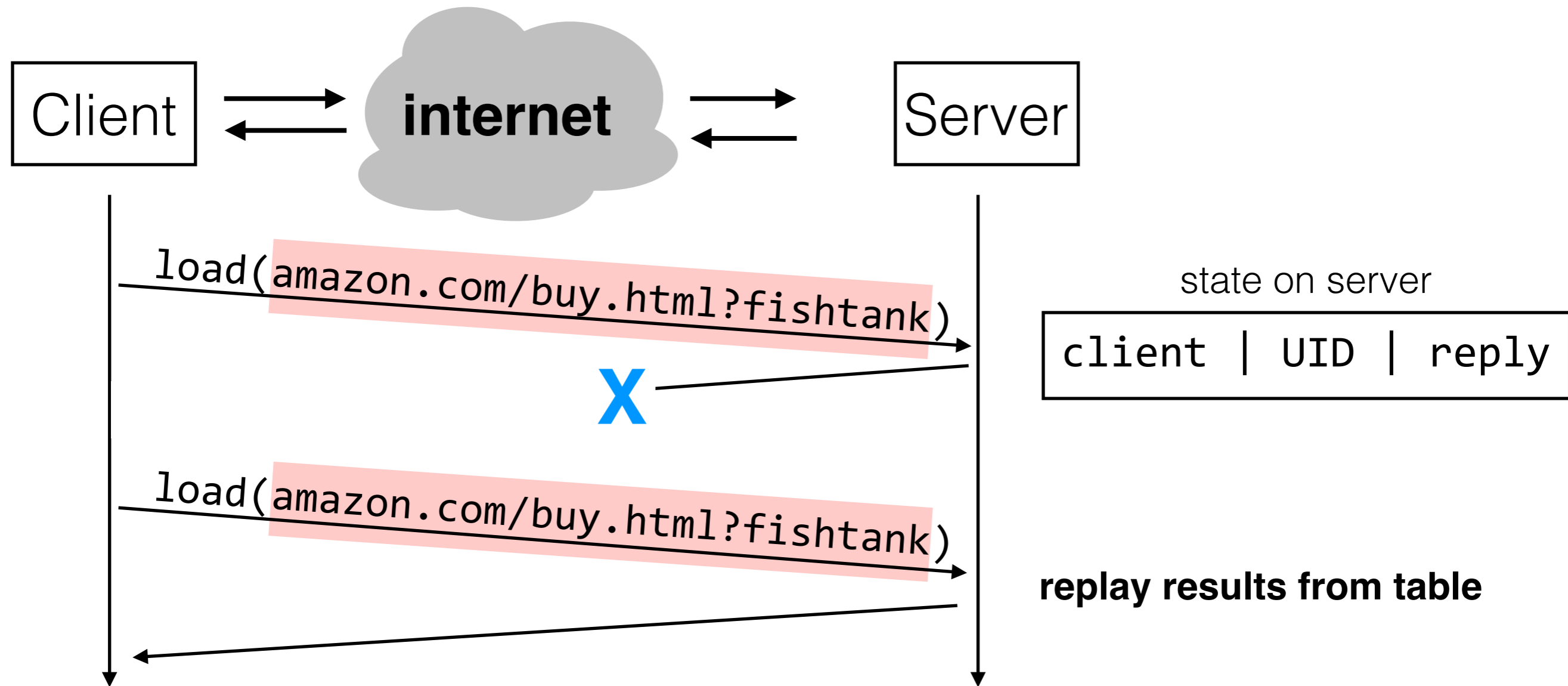


6.033 Spring 2015

Lecture #3

- Naming in systems
- Case study: DNS

Last Time: Enforced Modularity via Client/Server Model



Today: Naming

allows modules to interact

Naming

csail.mit.edu

katrina@csail.mit.edu

katrina

R0

main

WebBrowser

/mit/6.033/www/schedule.shtml

http://web.mit.edu/about

617-253-7341

128.30.2.121

hostname

email

username

x86 register name

function name

class name

path name

URL

phone number

IP Address

why use names?

Naming Schemes

1. Set of all possible **names**
2. Set of all possible **values**
3. **Look-up algorithm** to translate a name into a value (or set of values, or “none”)

Domain Name System

1. **names:** hostnames (`web.mit.edu`)

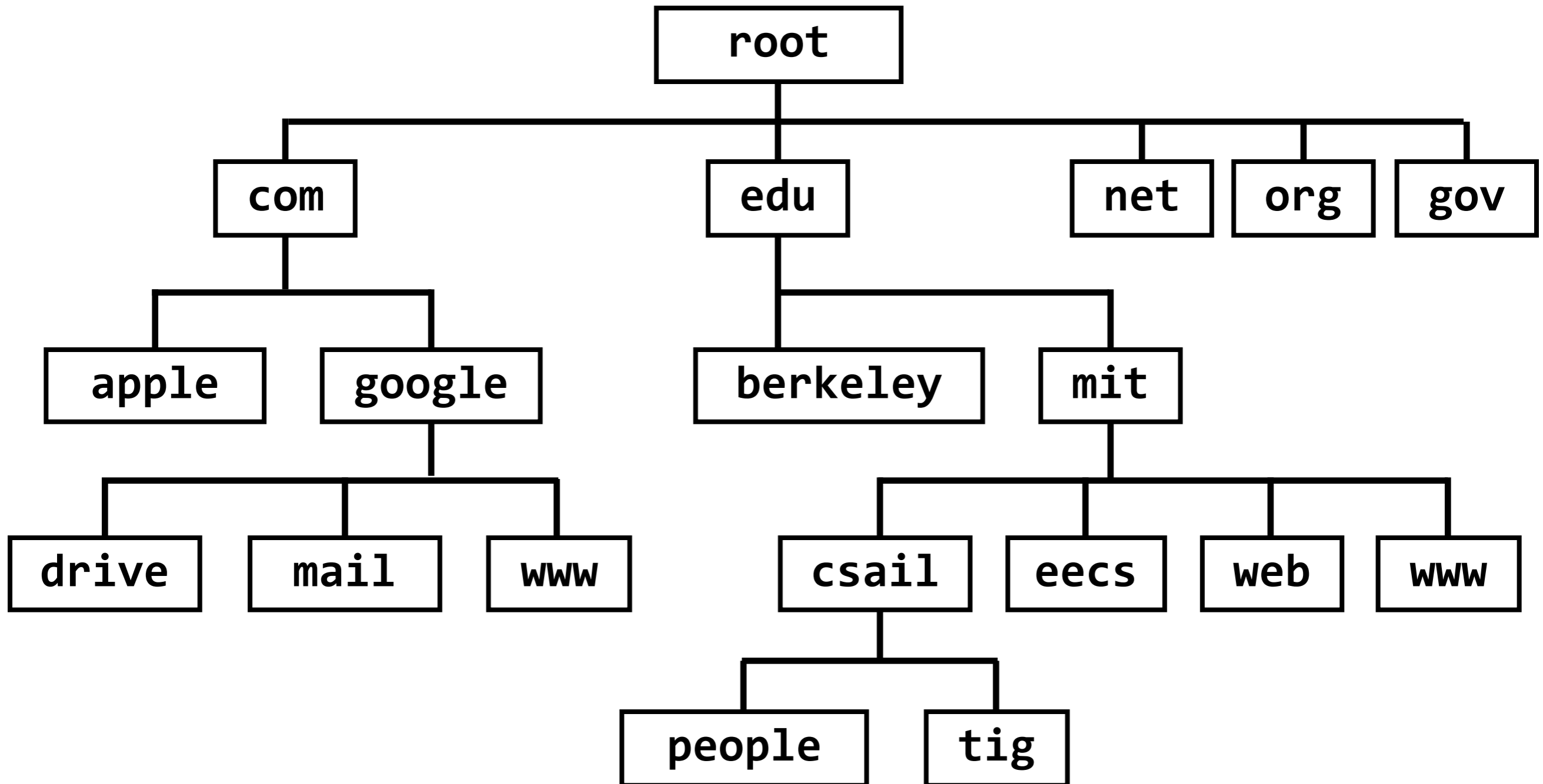
2. **values:** IP addresses (`18.9.22.69`)

IP addresses are imbued with location information: routers can send packets to an IP address, but not to a hostname

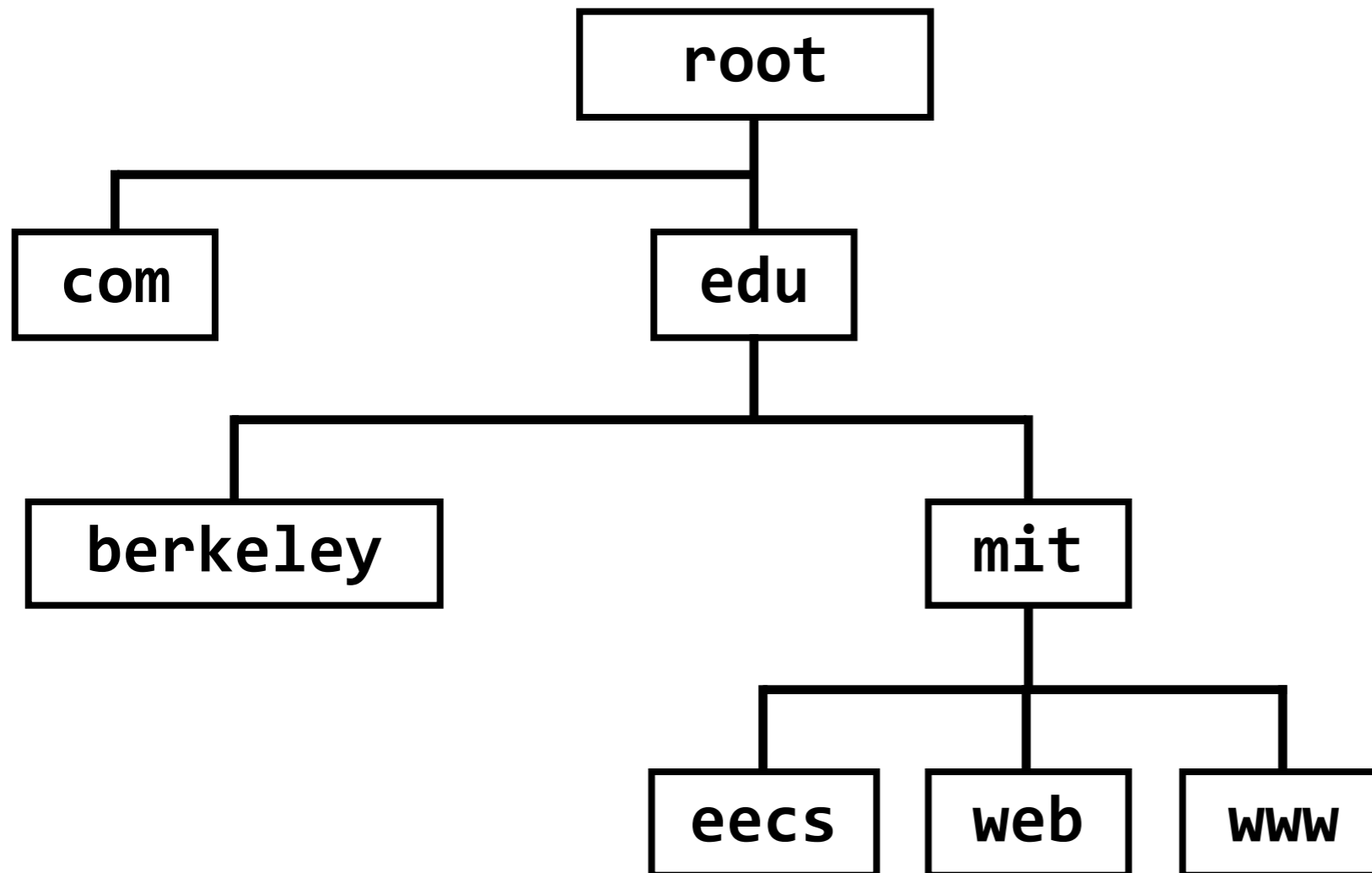
3. **look-up algorithm:** resolves a hostname to an IP address so that your machine knows where to send packets

DNS Hierarchy

(a partial view)



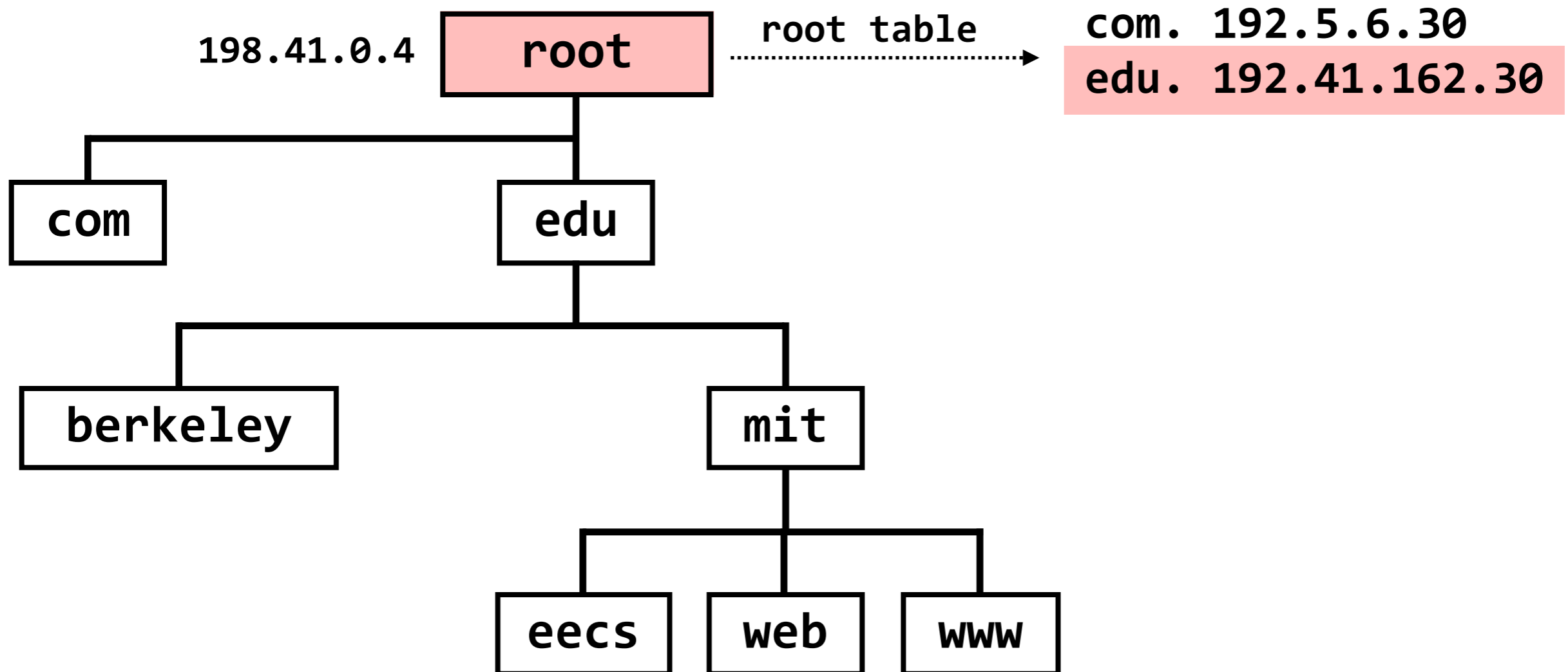
DNS Look-up



query to:

result:

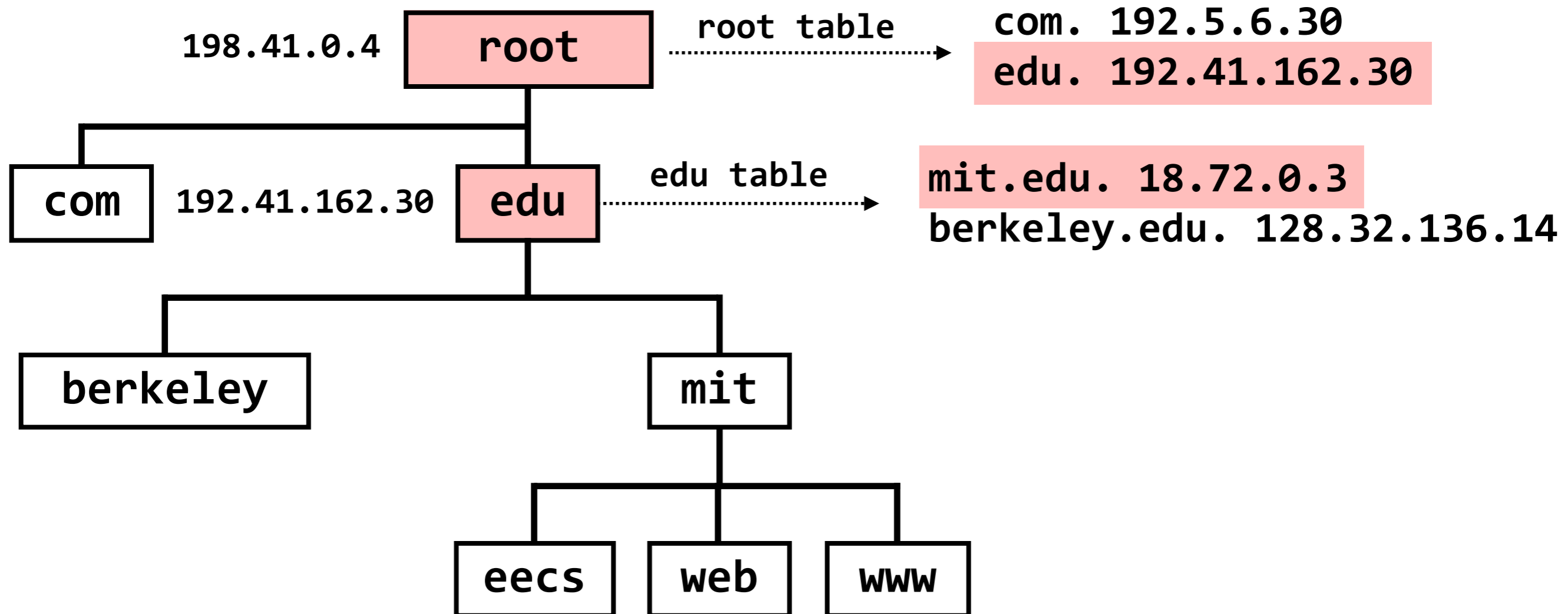
DNS Look-up



query to: 198.41.0.4

result: edu. 192.41.162.30

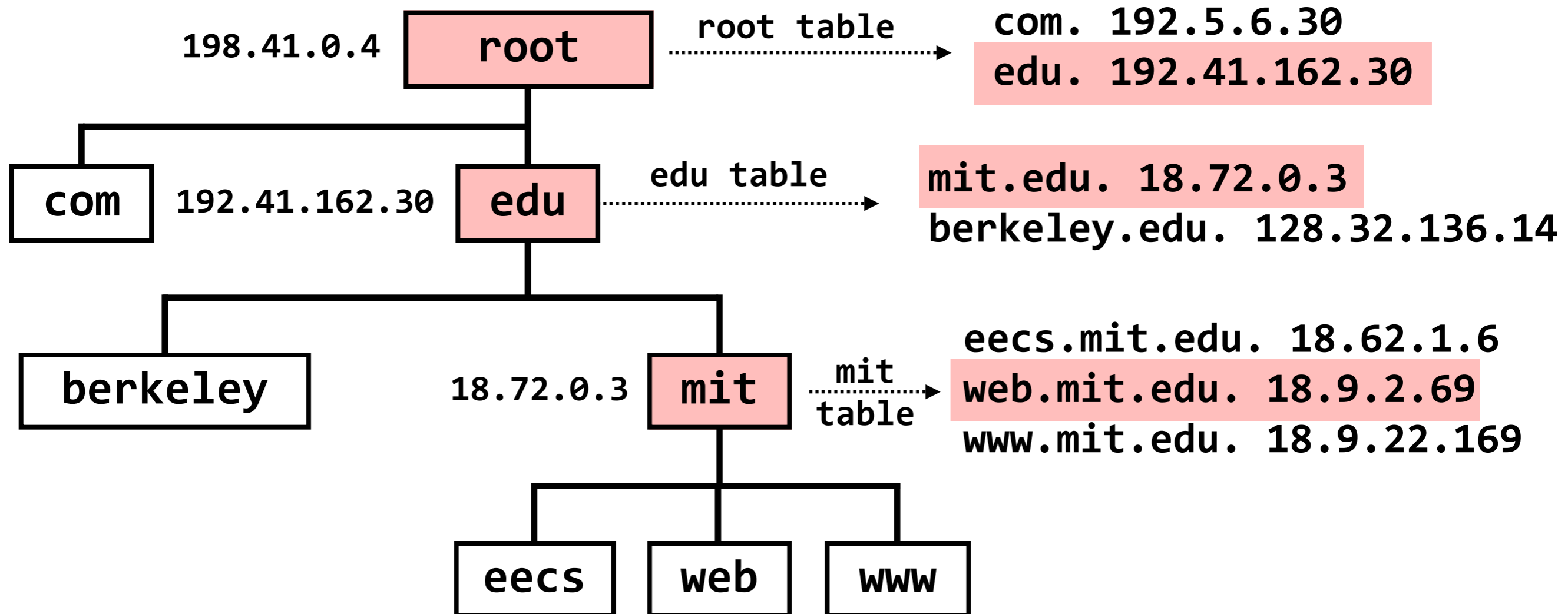
DNS Look-up



query to: 192.41.162.30

result: mit.edu. 18.72.0.3

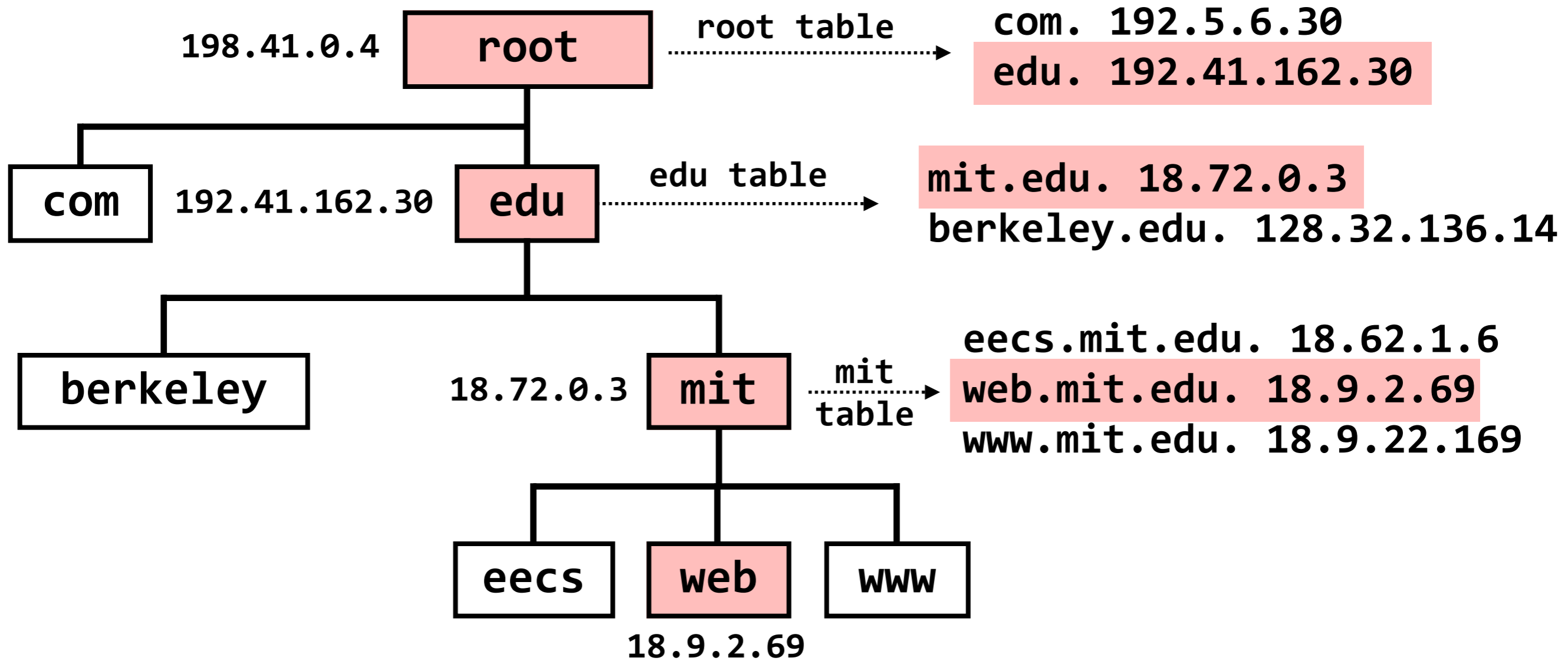
DNS Look-up



query to: 18.72.0.3

result: web.mit.edu. 18.9.2.69

DNS Look-up

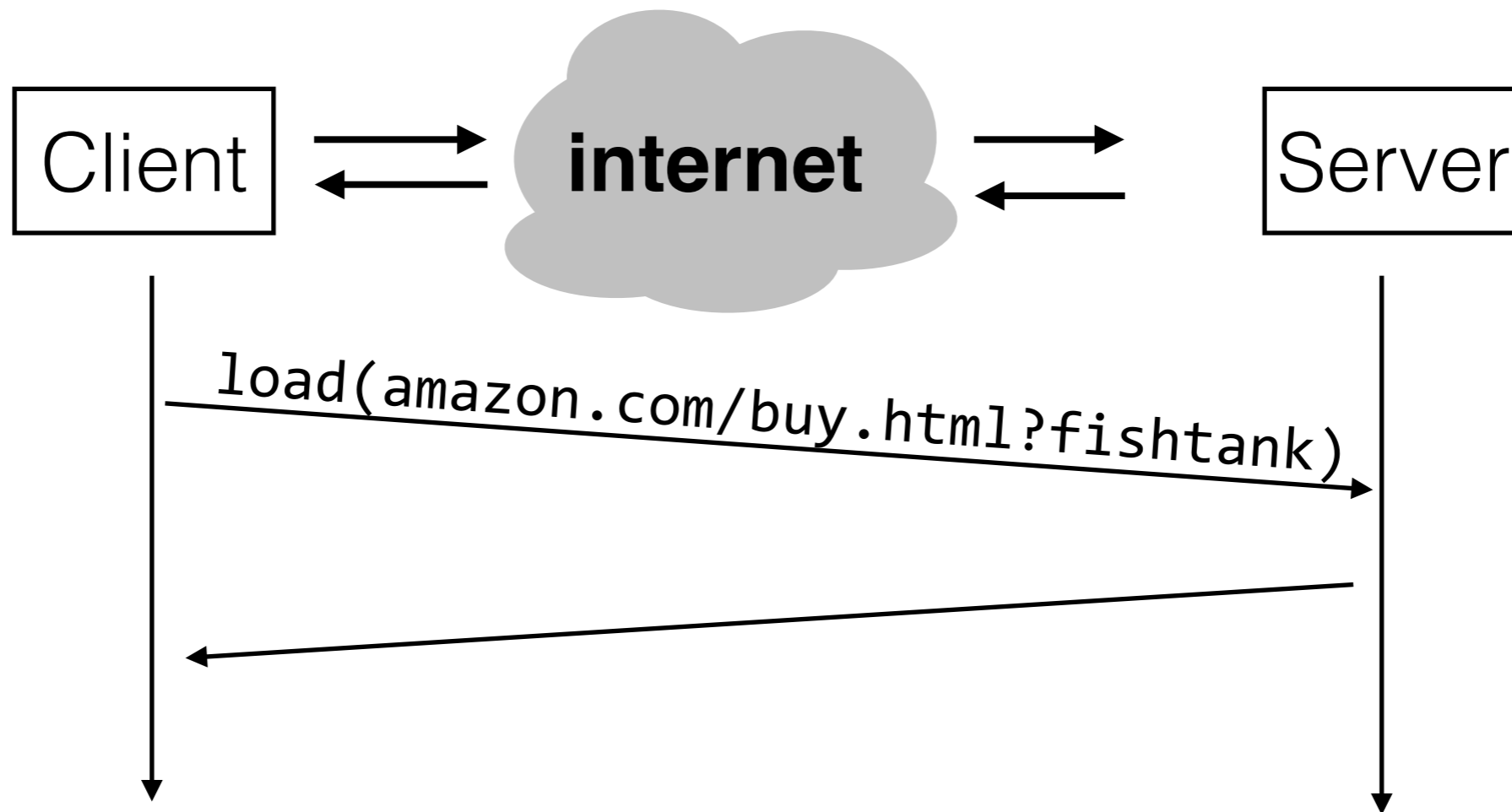


query to: 18.72.0.3

result: web.mit.edu. 18.9.2.69

- **Modularity** (previous lecture)
Modularity (and abstraction) limit complexity. One way to enforce modularity is to use a client/server design
- **Naming**
Naming is what allows modules — for example, a client and a server — to communicate; it is pervasive across systems
- **DNS**
DNS maps hostnames to IP addresses; its design is scalable and fault tolerant

Lingering Problem



what if we don't want our modules to be on entirely separate machines? how can we **enforce modularity on a single machine?**