

6.033: Intro to Computer Networks

Layering & Routing

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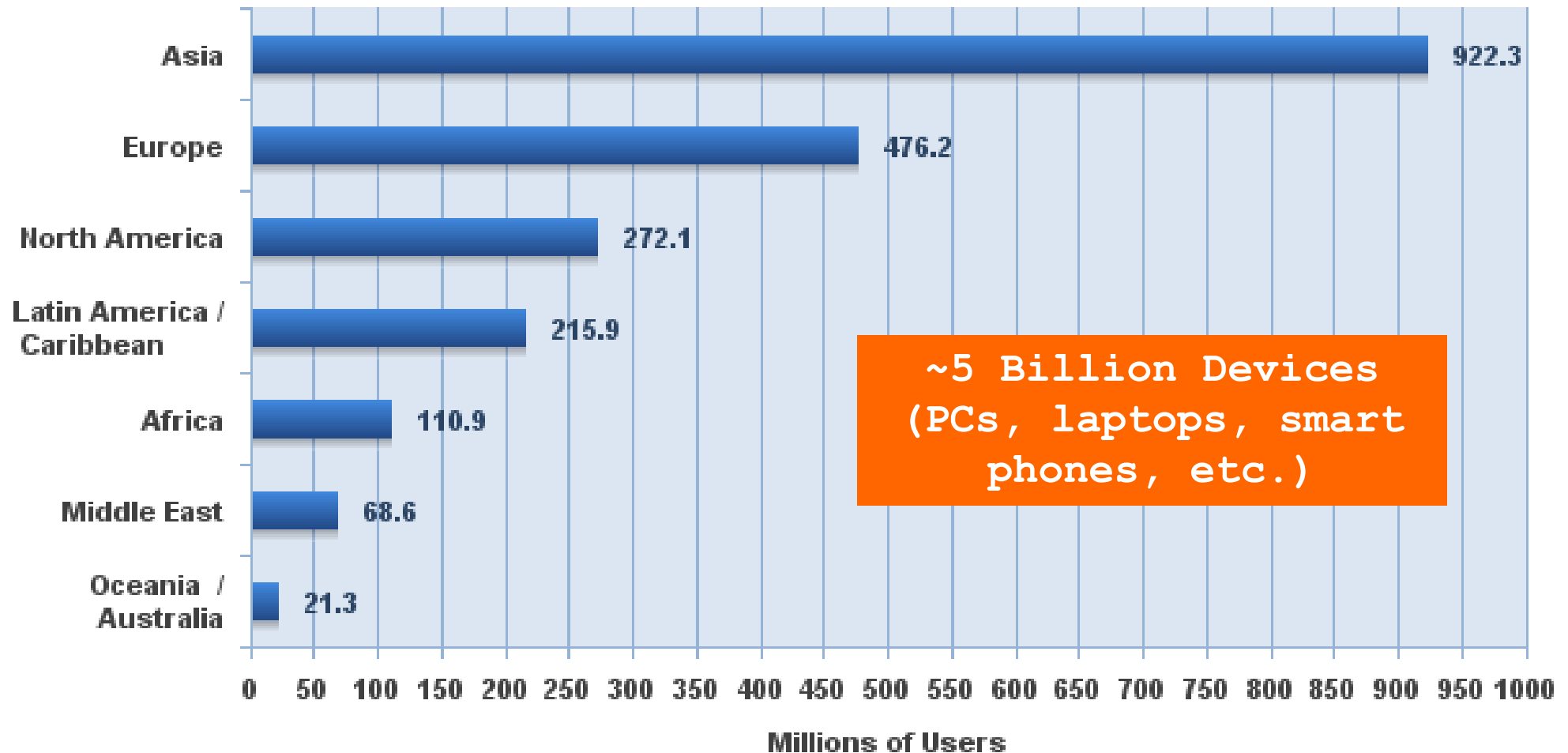
Some slides are contributed by N. McKewon, J. Rexford, I. Stoica



The Internet is an Exciting Place

Two Billion Internet Users

Internet Users in the World by Geographic Regions - 2011



Source: Internet World Stats - www.internetworldstats.com/stats.htm

Estimated Internet users are 2,095,006,005 on March 31, 2011

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The Internet is a Tense Place

Egypt blocks Internet access amid protests

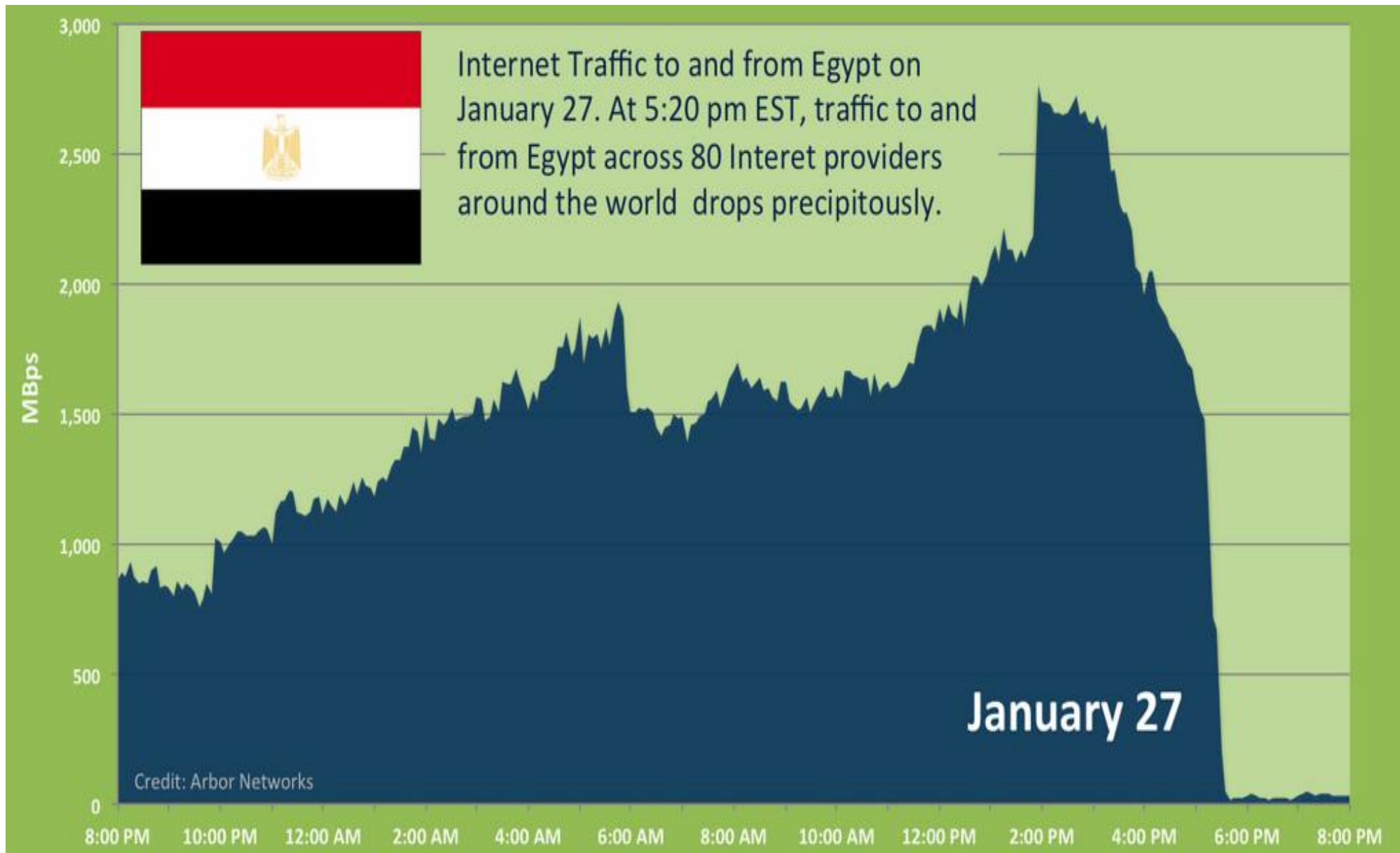
28 JANUARY 2011 Daniel Shane

Government orders telcos to block web access as protestors take to the streets

The Egyptian government has called on telecommunications providers in the country to block access to the Internet in response to widespread civil unrest.

Vodafone Egypt, one of the largest operators in the country not controlled by the state, today said it has disabled access following pressure from authorities.

Internet Traffic to/from Egypt



Stop Online Piracy Act (SOPA)



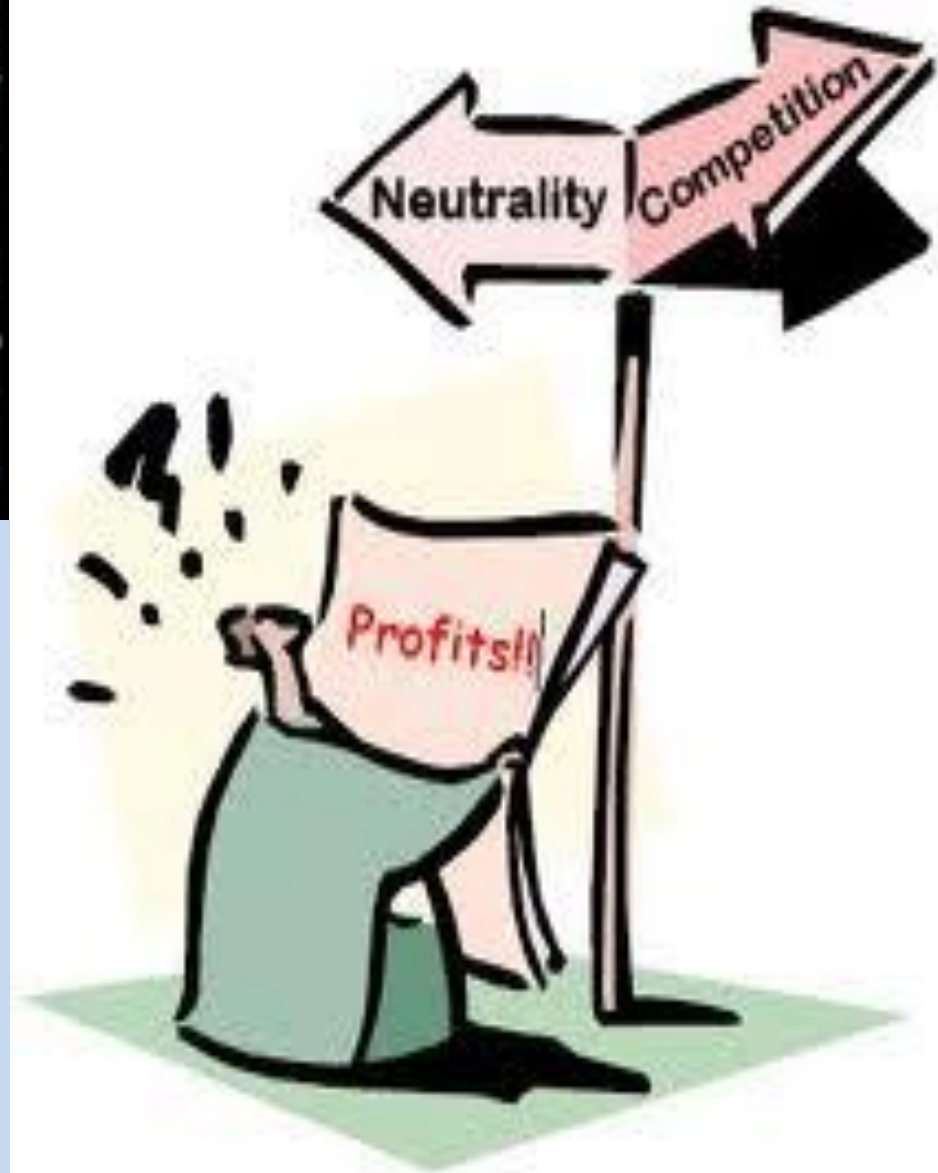
Network Neutrality

NET NEUTRALITY

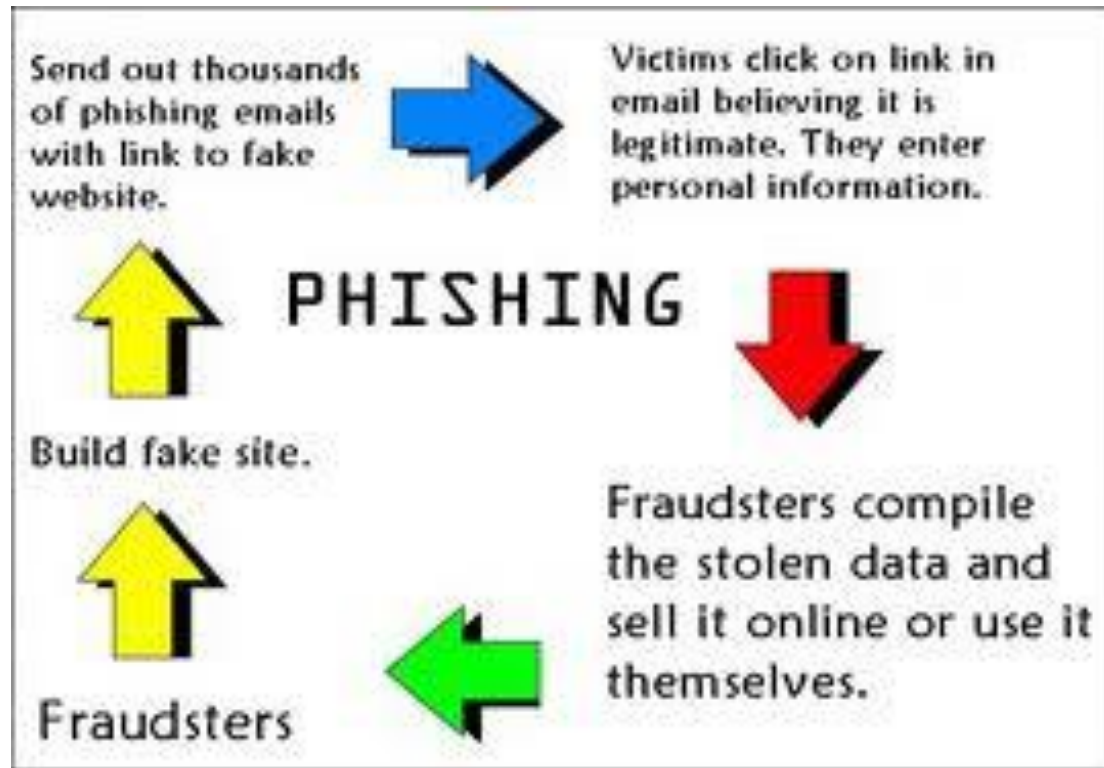
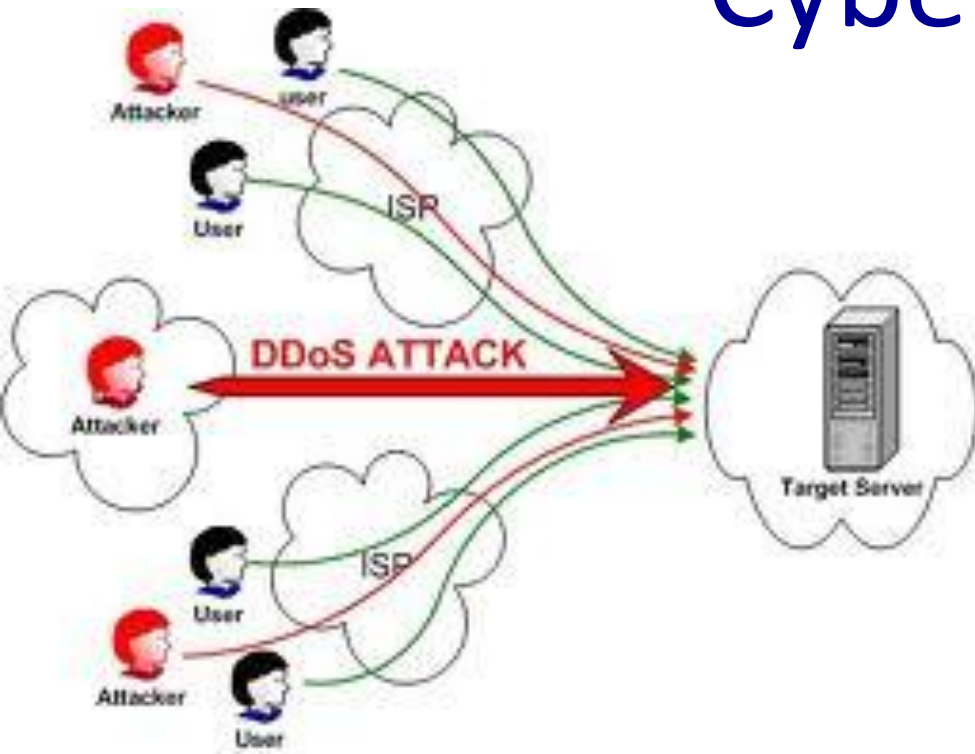
ALL BITS ARE CREATED EQUAL

FCC Rules Against Comcast P2P Throttling

The U.S. Federal Communications Commission has ordered Comcast to stop interfering with peer-to-peer traffic on its broadband network...



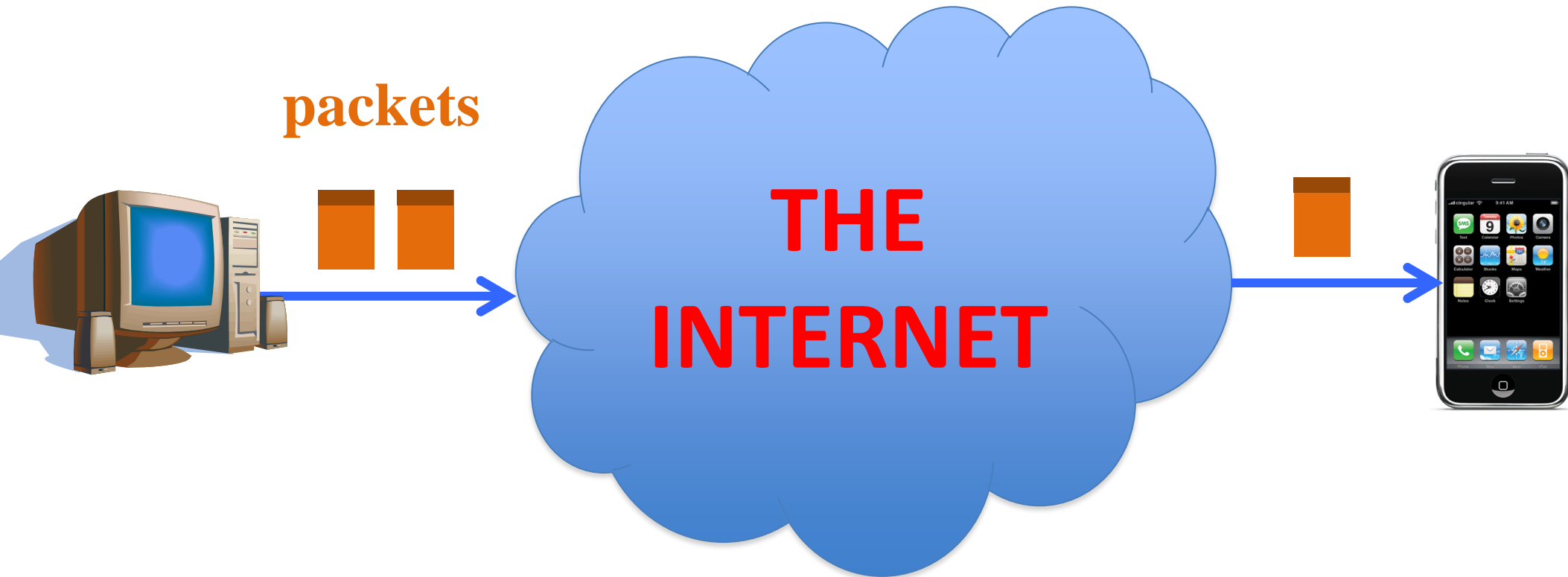
Cyber Attacks



What *is* the Internet?

3 guiding principles!

“Best-Effort Packet Delivery Service”



“Power at the Edge”

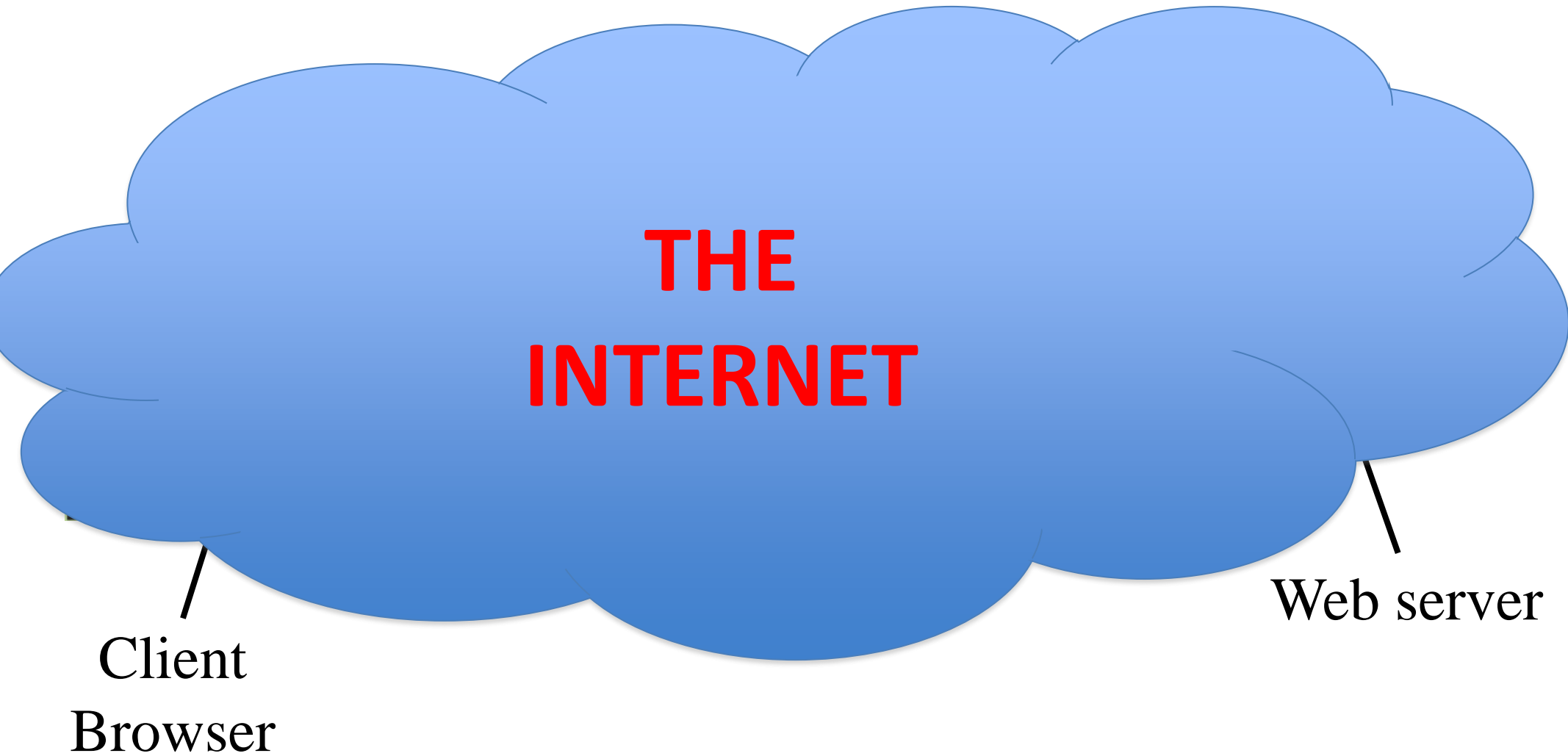
End-to-End Principle

Whenever possible, communications protocol operations should be defined to occur at the **end-points** of a communications system.

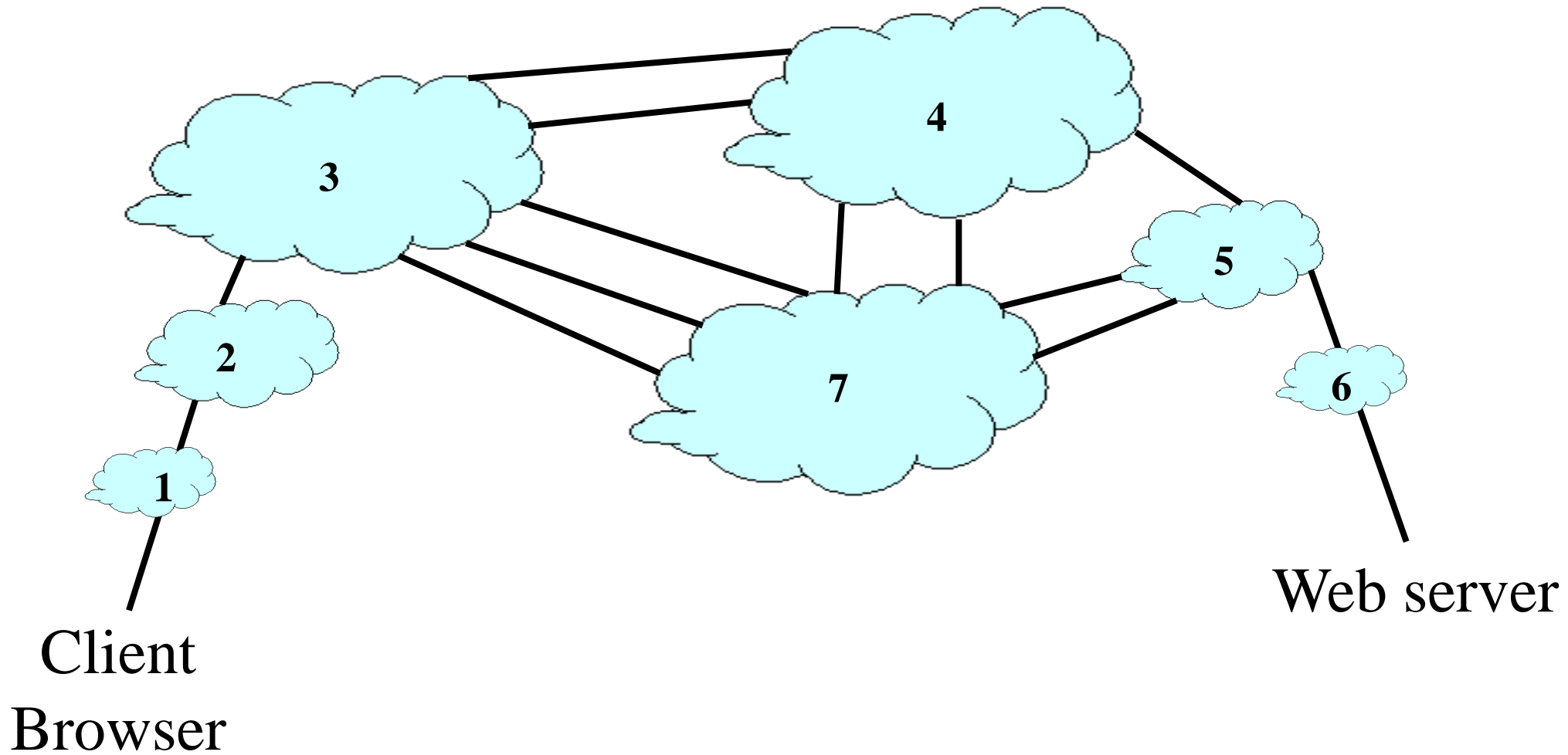
Keeps the network simple and scalable

Allows for easy introduction of new services at the edges

“A Network of Networks”



“A Network of Networks”

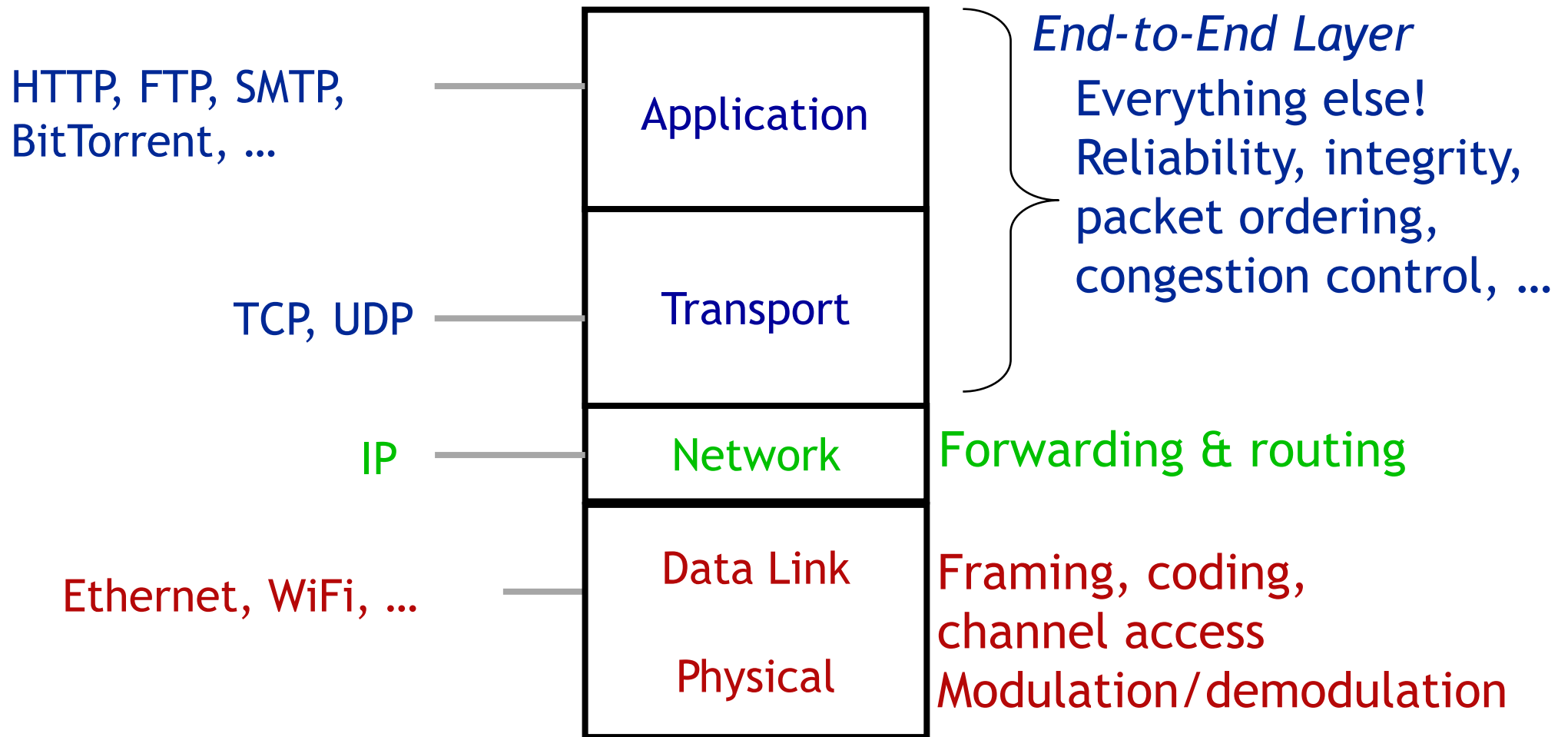


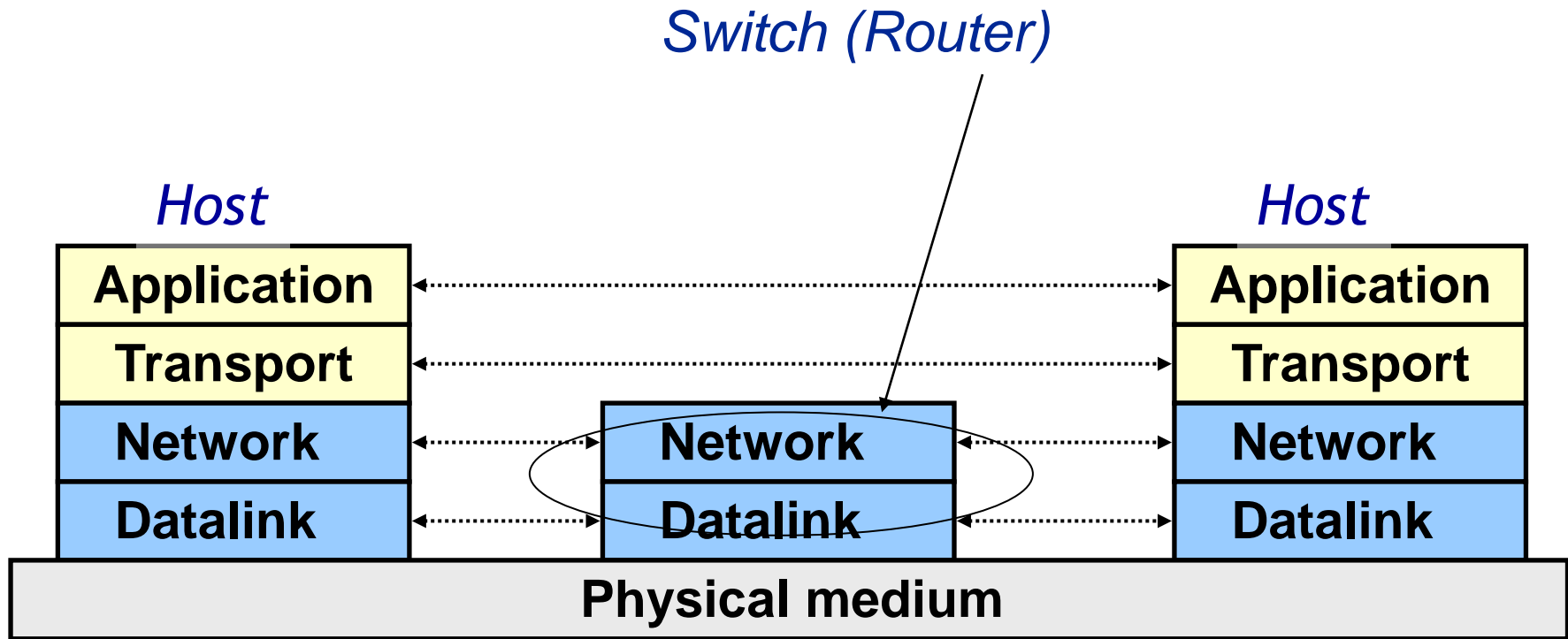
How the Internet is Organizing? Layering

Layering

- Layering is a particular form of abstraction
- The system is broken into a **vertical stack of functions/protocols**
- The service provided by one layer is based **solely** on the service provided by layer below
 - This is the “up/down” interface

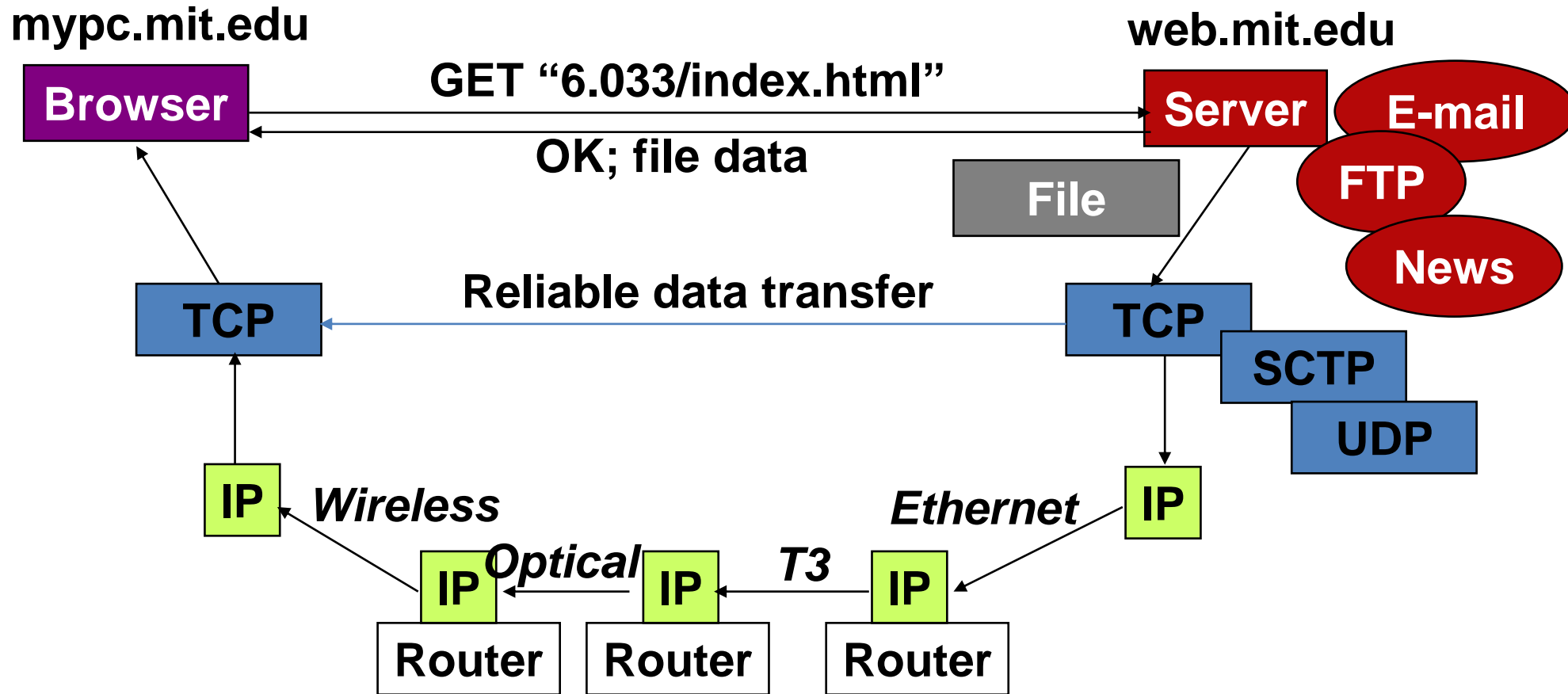
Layering in the Internet





- Link and network layers are implemented everywhere
- The end-to-end layer (i.e., transport and application) is implemented only at hosts

An Example



The Internet “Hourglass”

Applications

Web

FTP

Mail

News

Video

Audio

ping

napster

TCP

SCTP

UDP

ICMP

IP

Ethernet

802.11

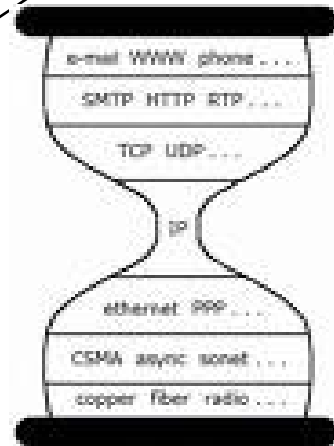
Power lines

ATM

Optical

Satellite

Bluetooth



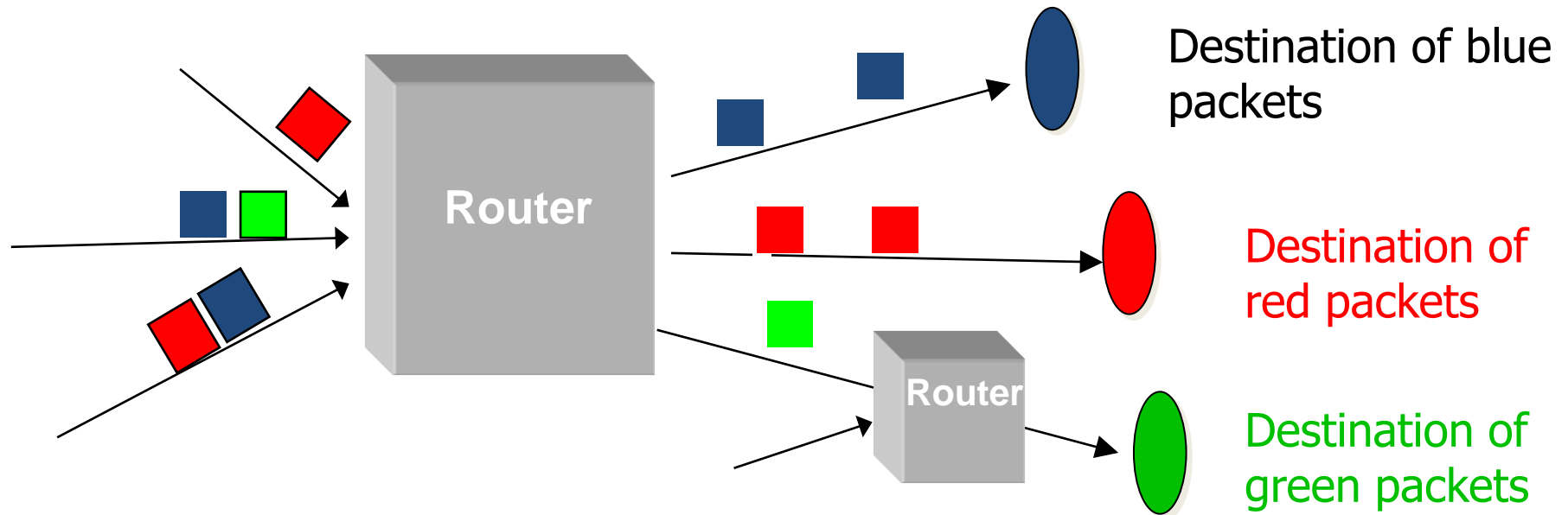
Link technologies

- Many applications, transports, and link protocols
- All use IP at the network layer: universal network layer

Network Layer

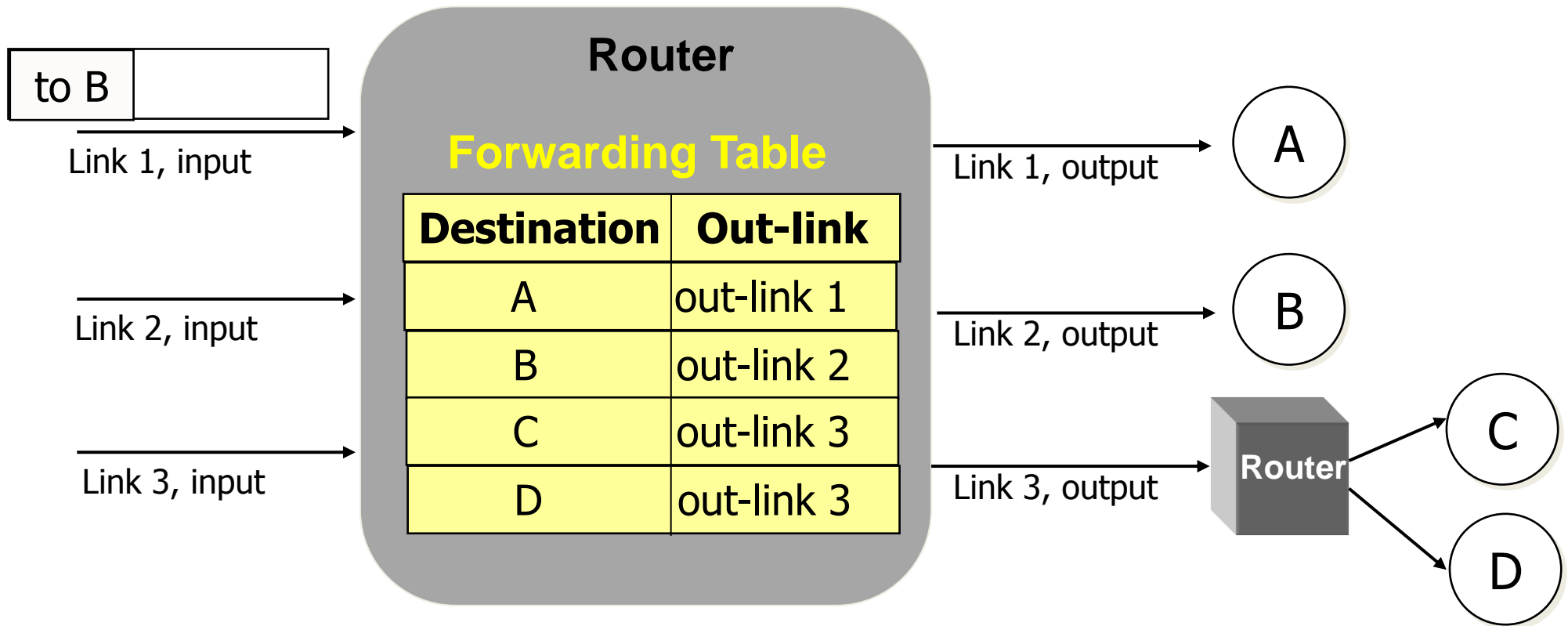
Routing (figuring out the routes)
&
Forwarding (sending the packets)

How Does a Router Forward the Packets?



- A router has input links and output links
- A router **sends** an input **packet** on the output link leading **toward the packet's destination node**
- A router does not care of who generated the packet

How does the router know which output link leads to a packet destination?



- Packet **header has the destination**
- Router **looks up the destination in its table** to find output link
- Table is built using a **routing protocol**

Basic Requirements of a Routing Protocol

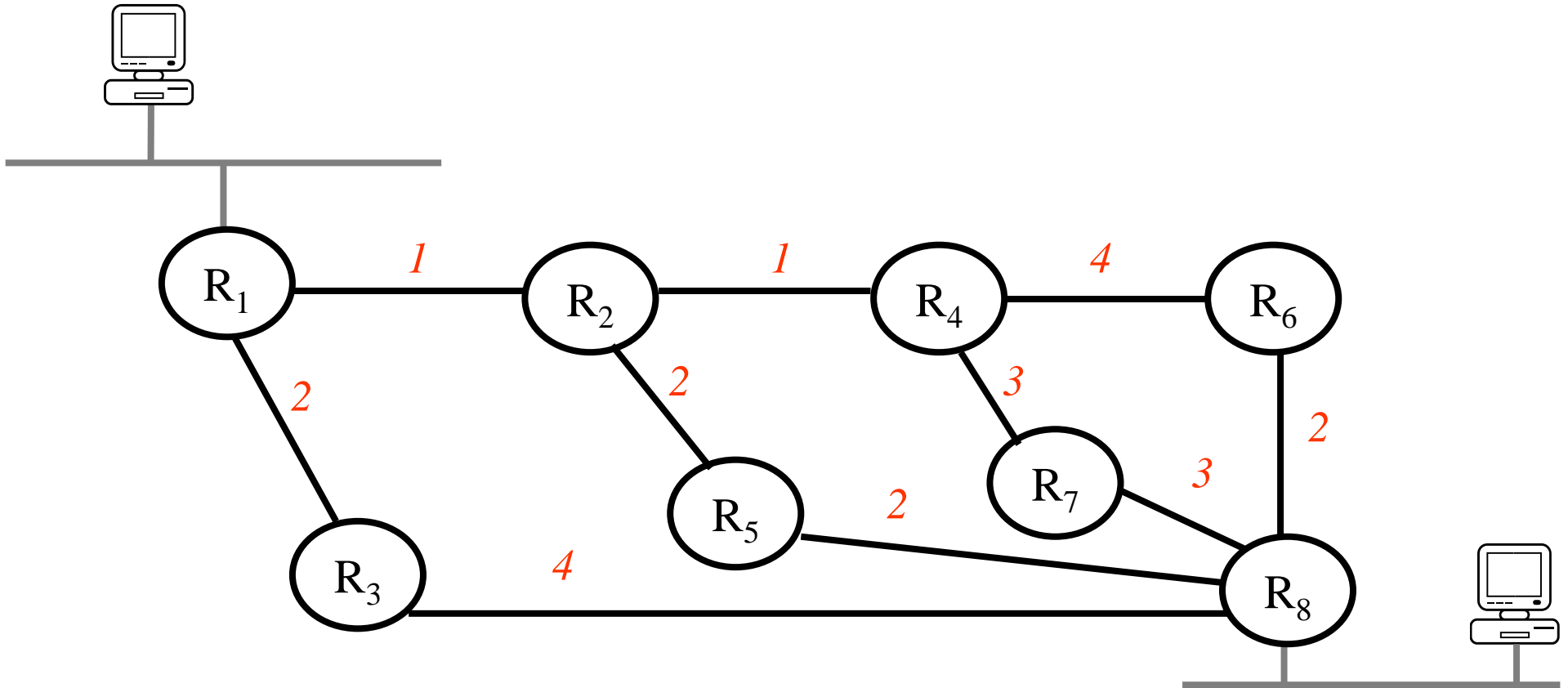
- Finds a path from source to destination
- Optimizes some metric (delay, cost, etc.)
- Has no (permanent) loops

Distance Vector Routing

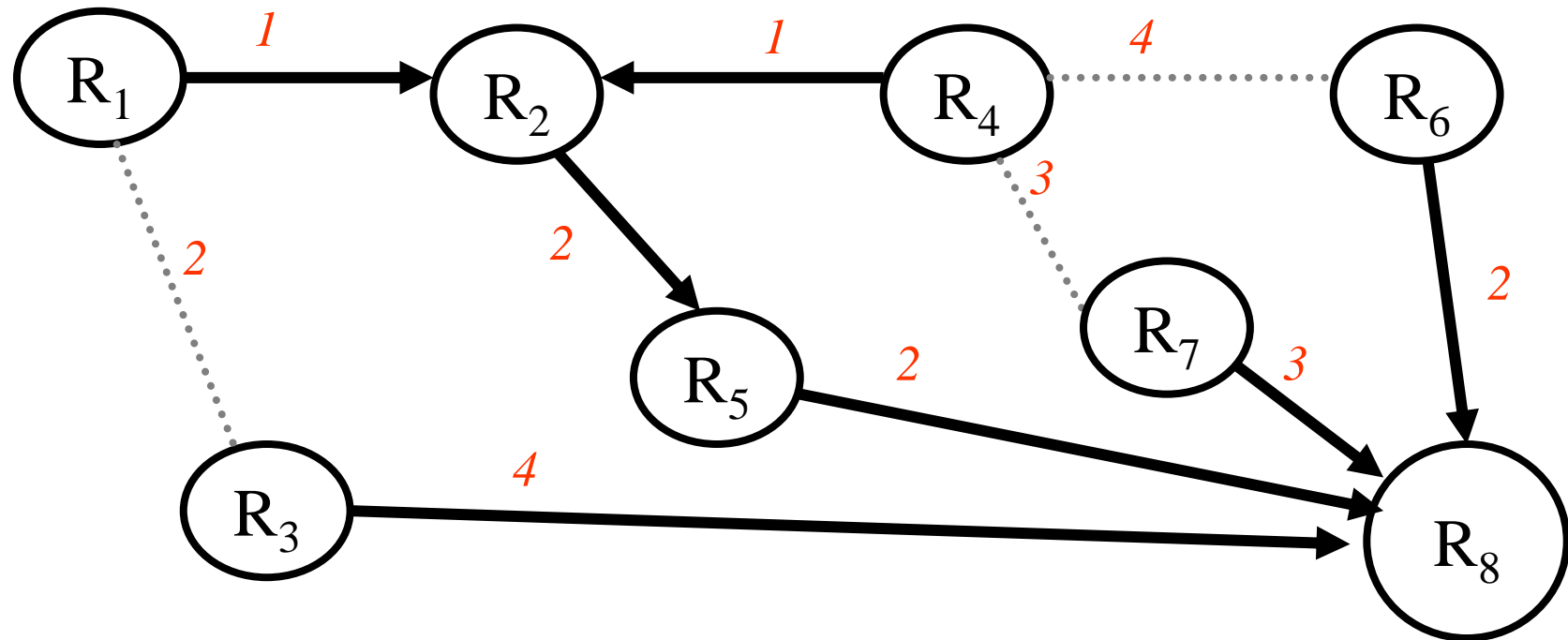
- Initialize
 - Distance to self is zero and next hop is self
 - Distance to anyone else is infinity
- Announce: Every T seconds
 - Tell neighbors distances to all destinations
- Update route to dst. upon message from j
 - Distance via $j = j$'s distance + weight of link to j
 - If distance via j is shorter than current distance, update routing table to go via j

Example

Objective: Determine the route from (R_1, \dots, R_7) to R_8 that minimizes the distance



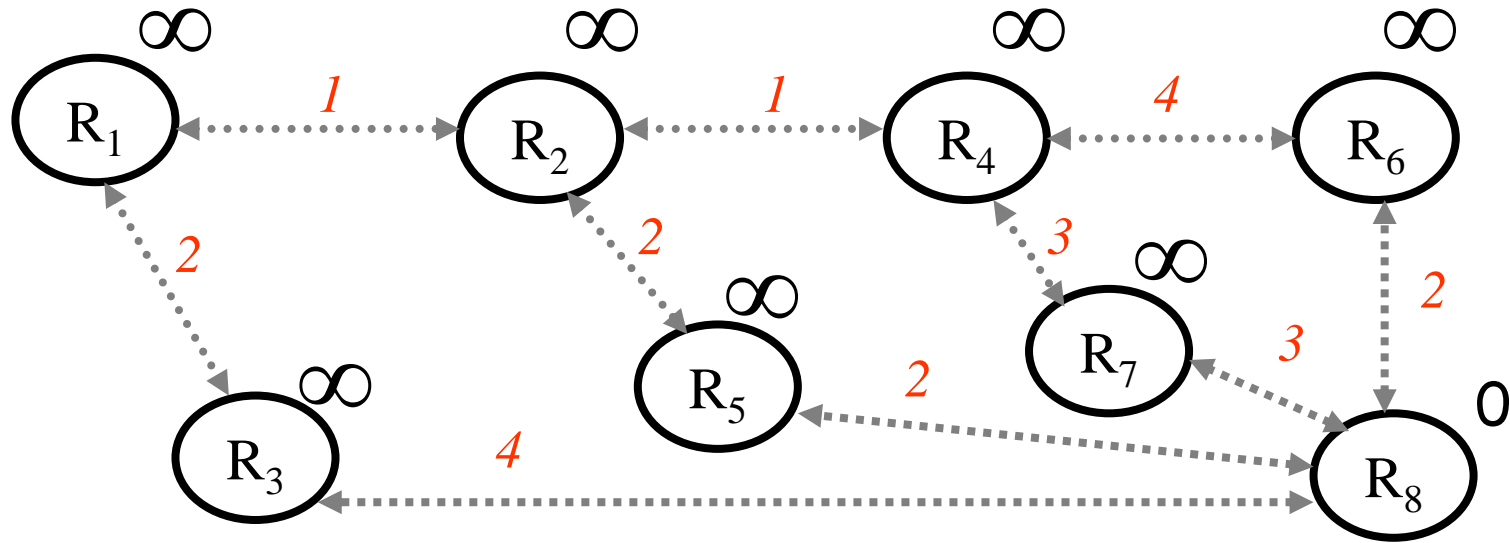
Solution is simple by inspection... (in this case)



The shortest paths from all sources to a destination (e.g., R_8) is the **spanning tree** routed at that destination.

Distance Vector Routing

Example

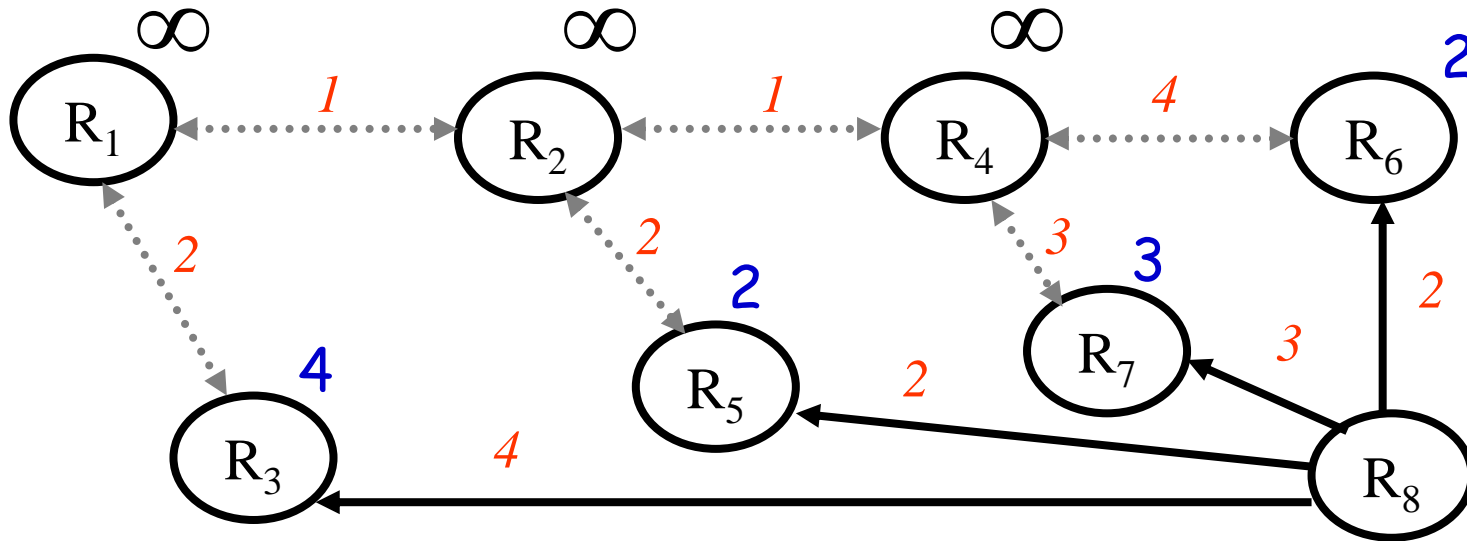


Initial State: All routers except R₈ set their route length to ∞ . R₈ sets its route length to 0.

Distance Vector Routing

Example

R ₁	Inf
R ₂	Inf
R ₃	4, R ₈
R ₄	Inf
R ₅	2, R ₈
R ₆	2, R ₈
R ₇	3, R ₈



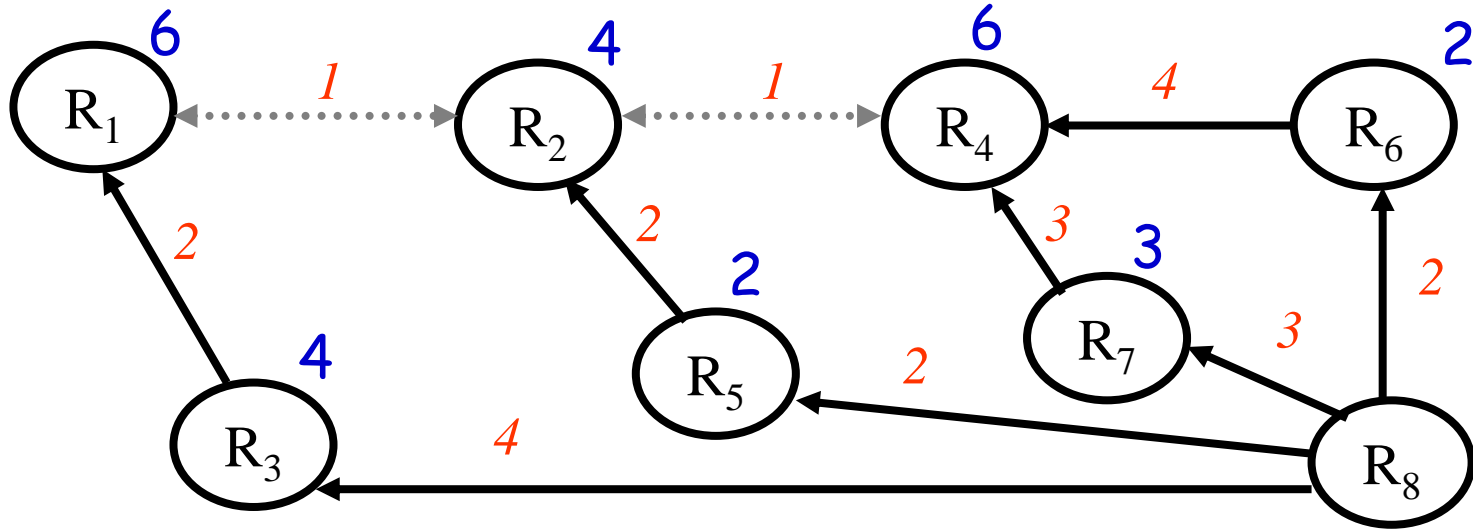
- ❖ Every T seconds, Router *i* tells its neighbors about its current lowest-cost path to R₈
- ❖ Each router updates its distance as $\min(\text{current distance}, \text{received distance} + \text{link weight})$

Note, routing tables have both the next-hop and the distance

Distance Vector Routing

Example

R ₁	6, R ₃
R ₂	4, R ₅
R ₃	4, R ₈
R ₄	6, R ₇
R ₅	2, R ₈
R ₆	2, R ₈
R ₇	3, R ₈

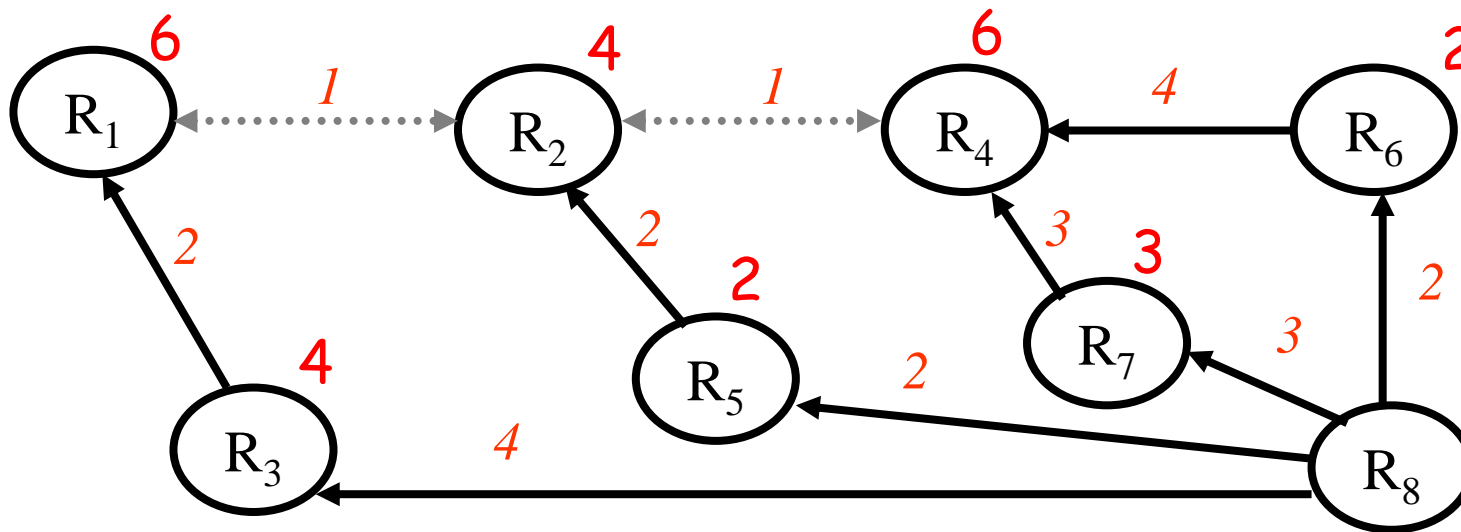


Repeat until no distance change

Distance Vector Routing

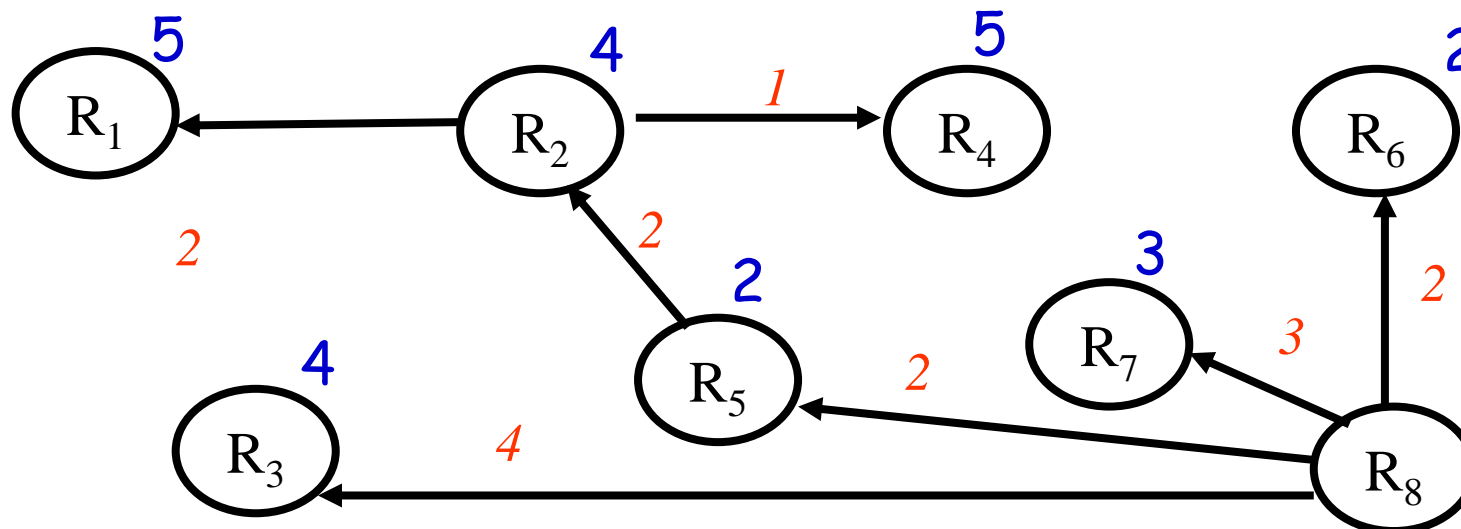
Example

R ₁	6, R3
R ₂	4, R5
R ₃	4, R ₈
R ₄	6, R7
R ₅	2, R ₈
R ₆	2, R ₈
R ₇	3, R ₈



Final Iteration

R ₁	5, R2
R ₂	4, R5
R ₃	4, R ₈
R ₄	5, R2
R ₅	2, R ₈
R ₆	2, R ₈
R ₇	3, R ₈



Summary

- Internet architecture is based on layering
 - E2E Layer
 - Network Layer
 - Link Layer
- Job of Network Layer is Routing & Forwarding
 - Routers build routing tables using routing protocol
 - Routers forward packets based on the packet's header and the routing table