

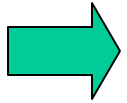
Chapter 7

An Introduction to Computer Networks

Prof. Dina Katabi

Some slides are from lectures by Nick Mckeown, Ion Stoica, Frans Kaashoek, Hari Balakrishnan, and Sam Madden

Chapter Outline



Introduction (slides and 7.A)

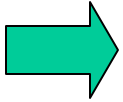
Layered Architecture (slides and 7.B & 7.D)

Routing (slides and 7.D)

Reliable Transmission & Flow Control (slides and read 7.E)

Congestion Control (slides and read 7.F)

This Lecture



What is a network?

Sharing the infrastructure

Circuit switching

Packet switching

Best Effort Service

Analogy: the mail system

Internet's Best Effort Service

Networks

Why they are interesting?

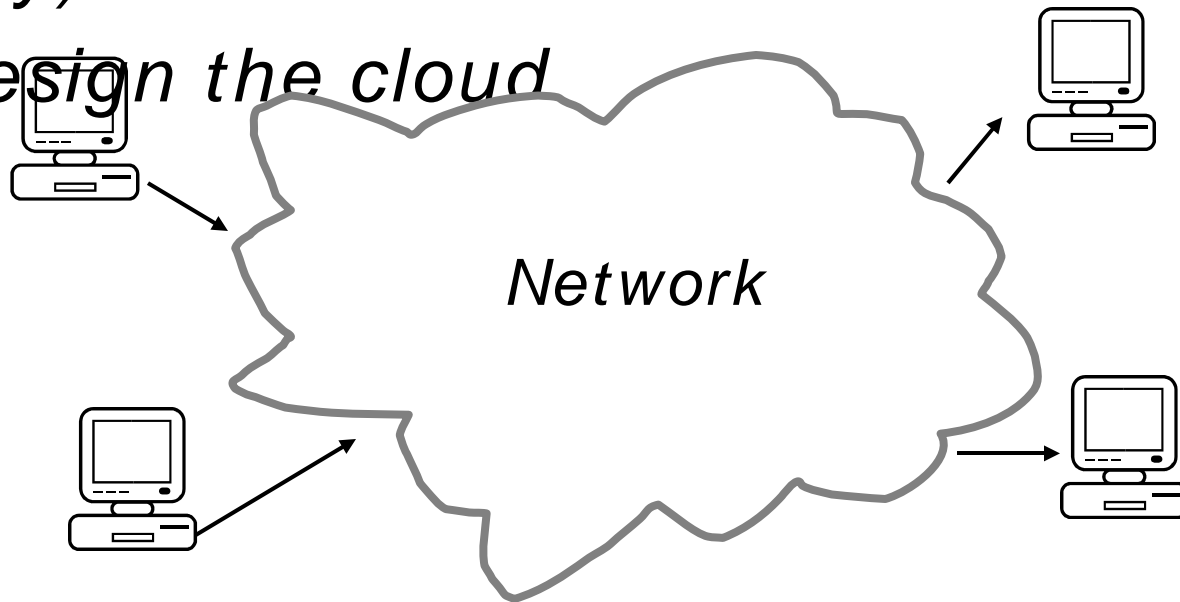
Overcome geographic limits

Access remote data

Separate clients and server

Goal: Universal Communication (any to any)

Design the cloud



Connectivity



Link

DSL, T1, T3, ...

Characterized by

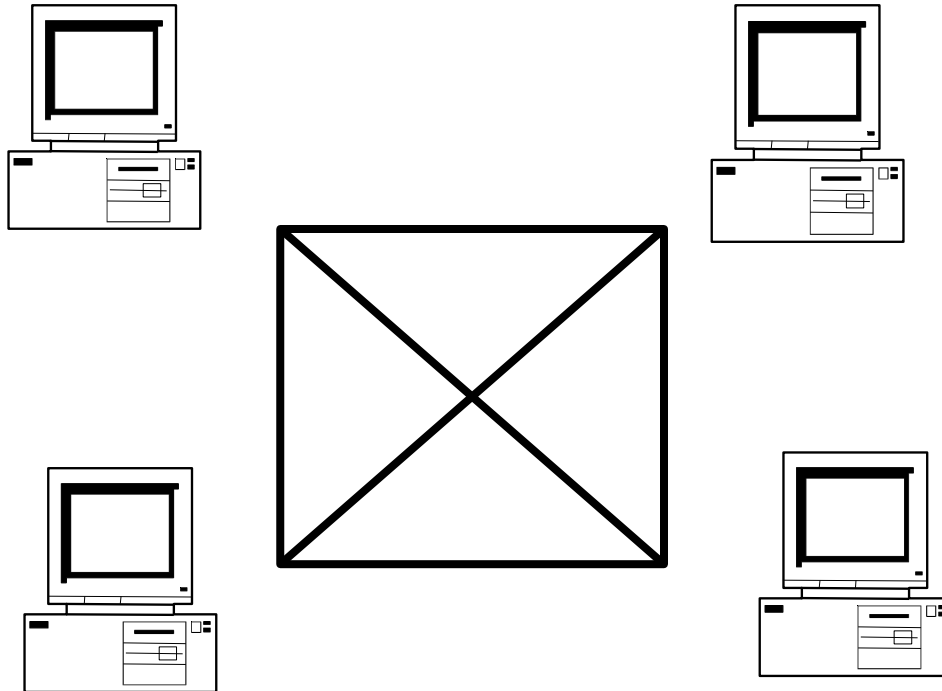
Capacity or bit-rate (1.5 Mb/s, 100Mb/s, ...)

Propagation delay (10us, 10ms, 100ms, ..)

Transfer time on a link = # bit/bit-rate + propagation delay

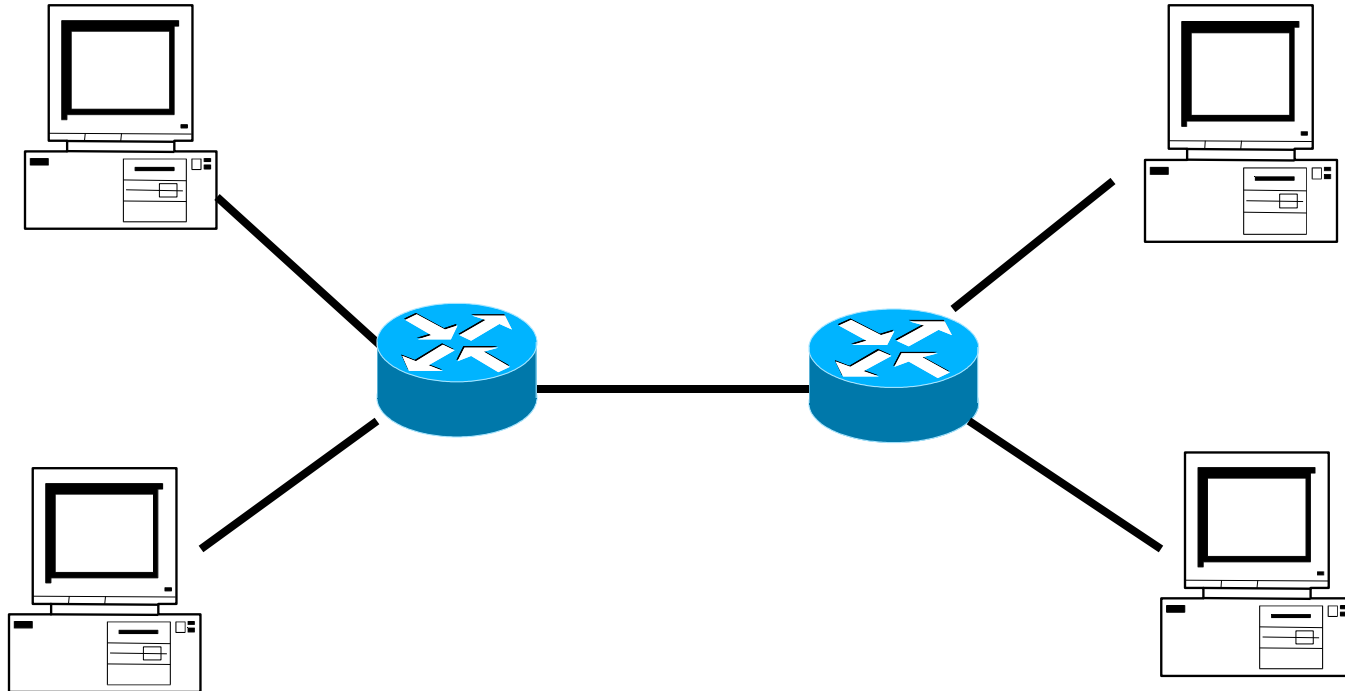
Connectivity

A mesh requires N^2 links too costly



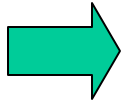
We Have to Share the Infrastructure

Intermediate nodes called switches or routers allow the hosts to share the infrastructure



This Lecture

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Two ways to share

Circuit switching (isochronous)

Packet switching (asynchronous)

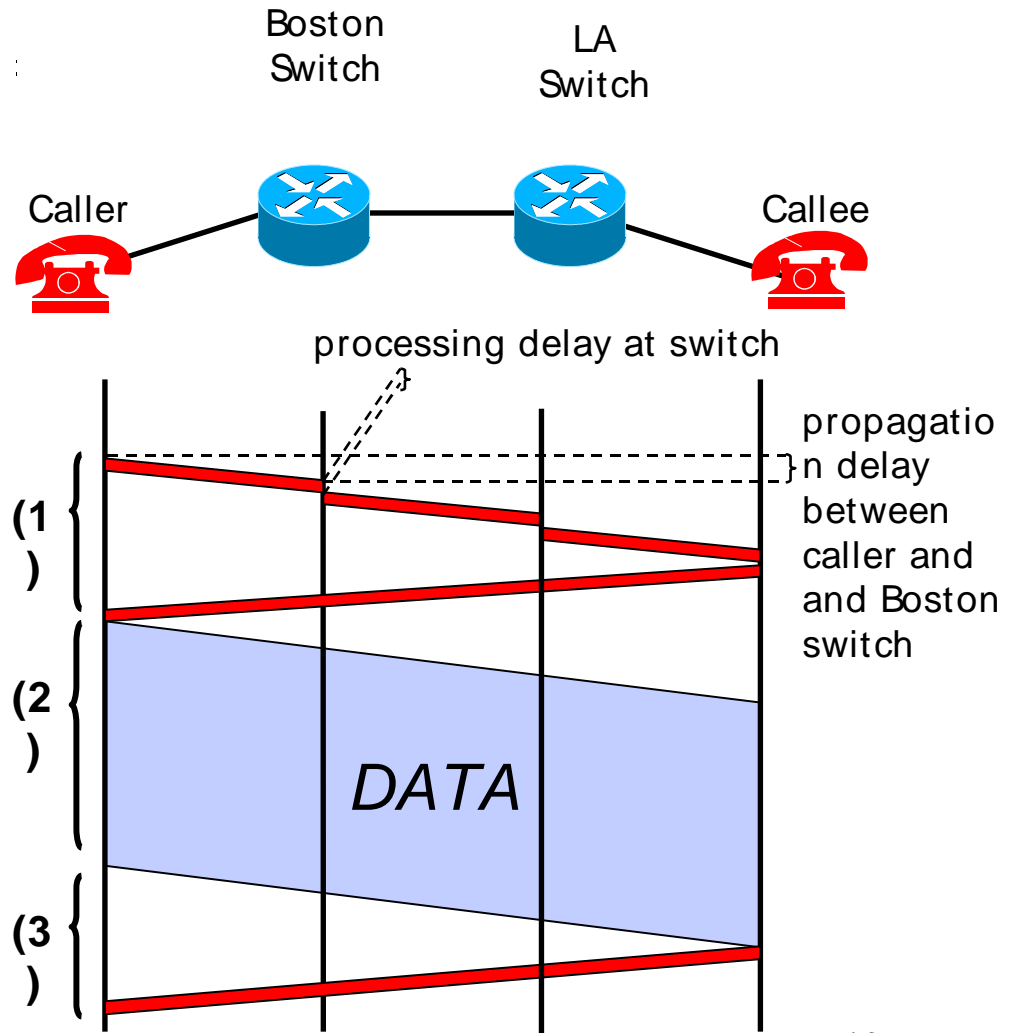
Circuit Switching

It's the method used by the telephone network

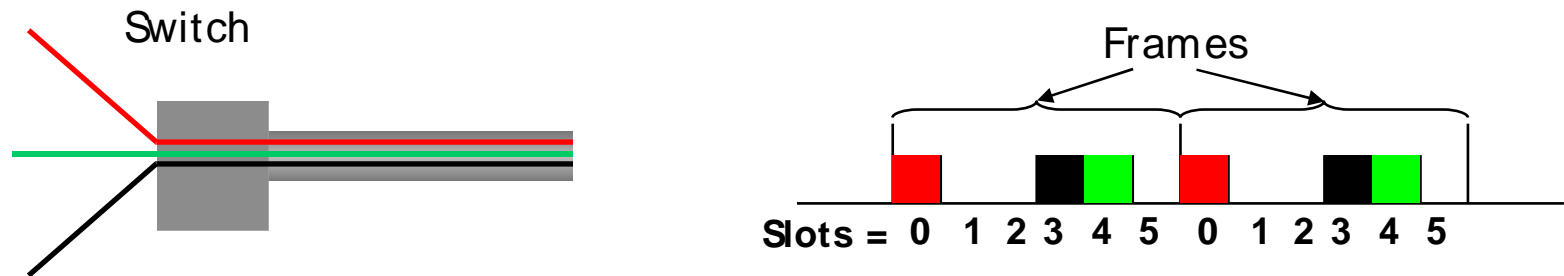
A call has three phases:

1. *Establish circuit from end-to-end ("dialing"),*
2. *Communicate,*
3. *Close circuit ("tear down").*

If circuit not available: "busy signal"



Circuit Switching: Multiplexing/Demultiplexing



One way for sharing a circuit is TDM:

*Time divided into frames and frames divided into slots
Relative slot position inside a frame determines which
conversation the data belongs to*

E.g., slot 0 belongs to the red conversation

Need synchronization between sender and receiver

Circuit Switching

Assume link capacity is C bits/sec

Each communication requires R bits/sec

slots = C/R

Maximum number of concurrent communications is C/R

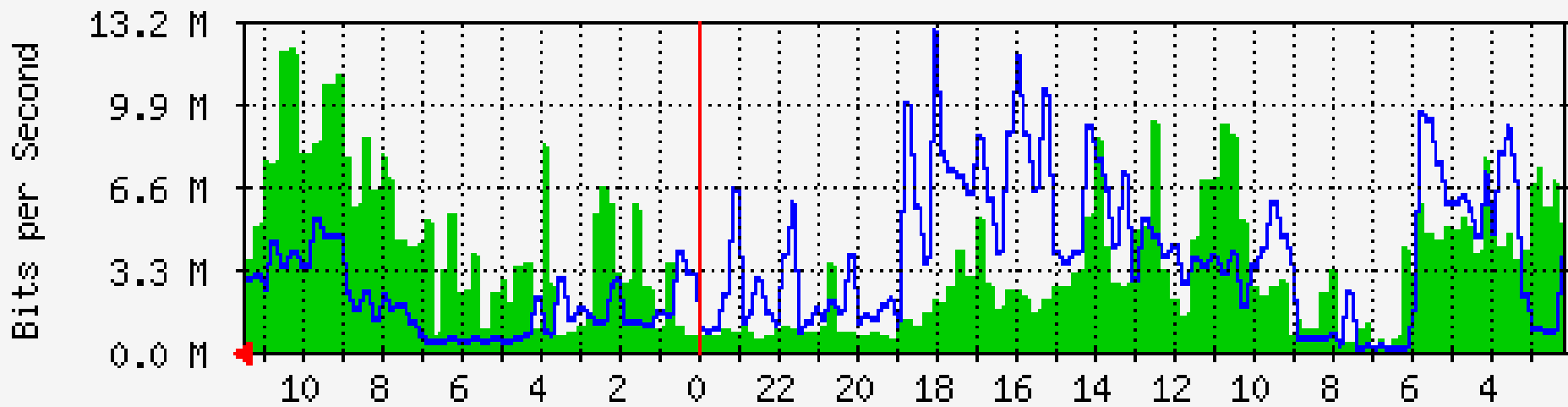
What happens if we have more than C/R communications?

What happens if the a communication sends less/more than R bits/sec?

Design is unsuitable for computer networks where transfers have variable rate (bursty)

Internet Traffic Is Bursty

Daily traffic at an MIT-CSAIL router



Max In: 12.2 Mb/s

Avg. In: 2.5 Mb/s

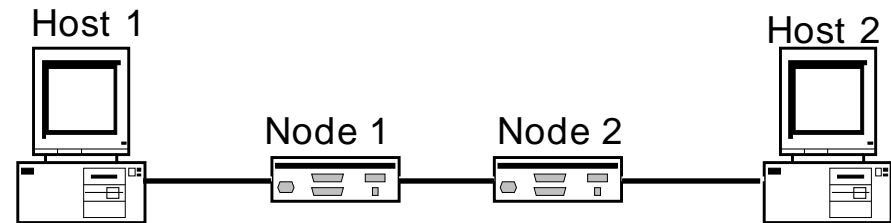
Max Out: 12.8 Mb/s

Avg. Out: 3.4 Mb/s

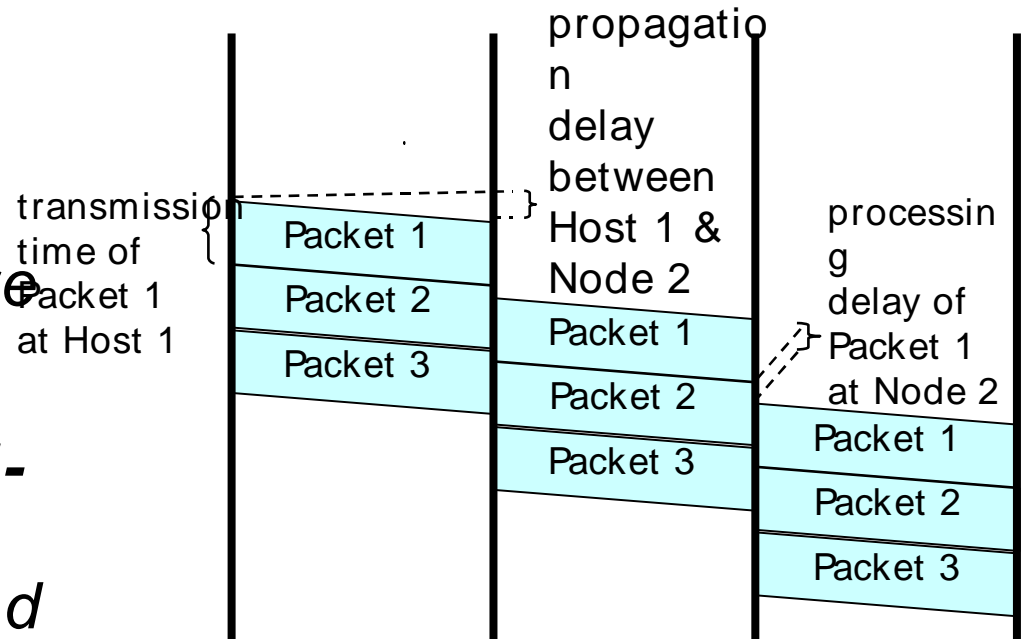
Packet Switching

Used in the Internet

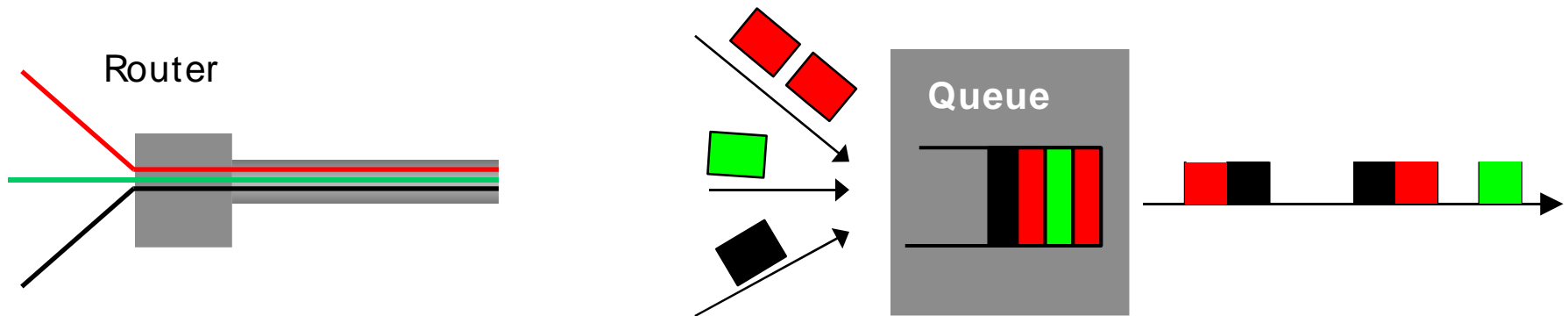
*Data is sent in **Packets**
(header contains control
info, e.g., source and
destination addresses)*



*Per-packet routing
At each node the entire
packet is received,
stored, and then
forwarded (**store-and-
forward networks**)
No capacity is allocated*



Packet Switching: Multiplexing/Demultiplexing



Multiplex using a queue

Routers need memory/buffer

Demultiplex using information in packet header

Header has destination

Router has a routing table that contains information about which link to use to reach a destination

Queues introduce

Variable Delay

*Delay = Queuing delay + propagation delay +
transmission delay + processing delay*

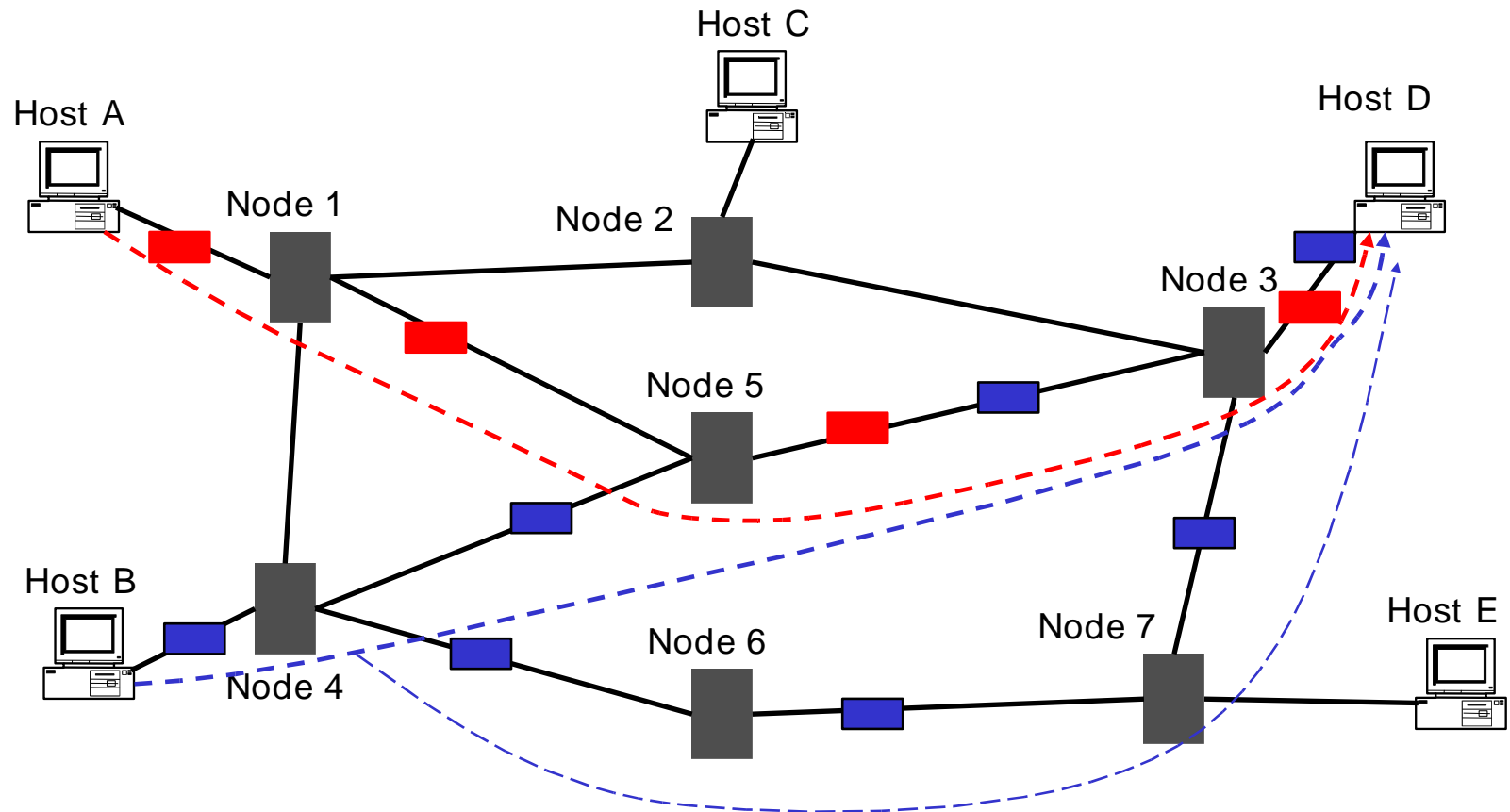
Losses

*When packets arrive to a full queue/buffer they are
dropped*

Packet switching also show

reordering

Packets in a flow may not follow the same path (depends on routing as we will see later) packets may be reordered



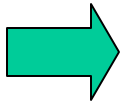
This Lecture

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Packet switching

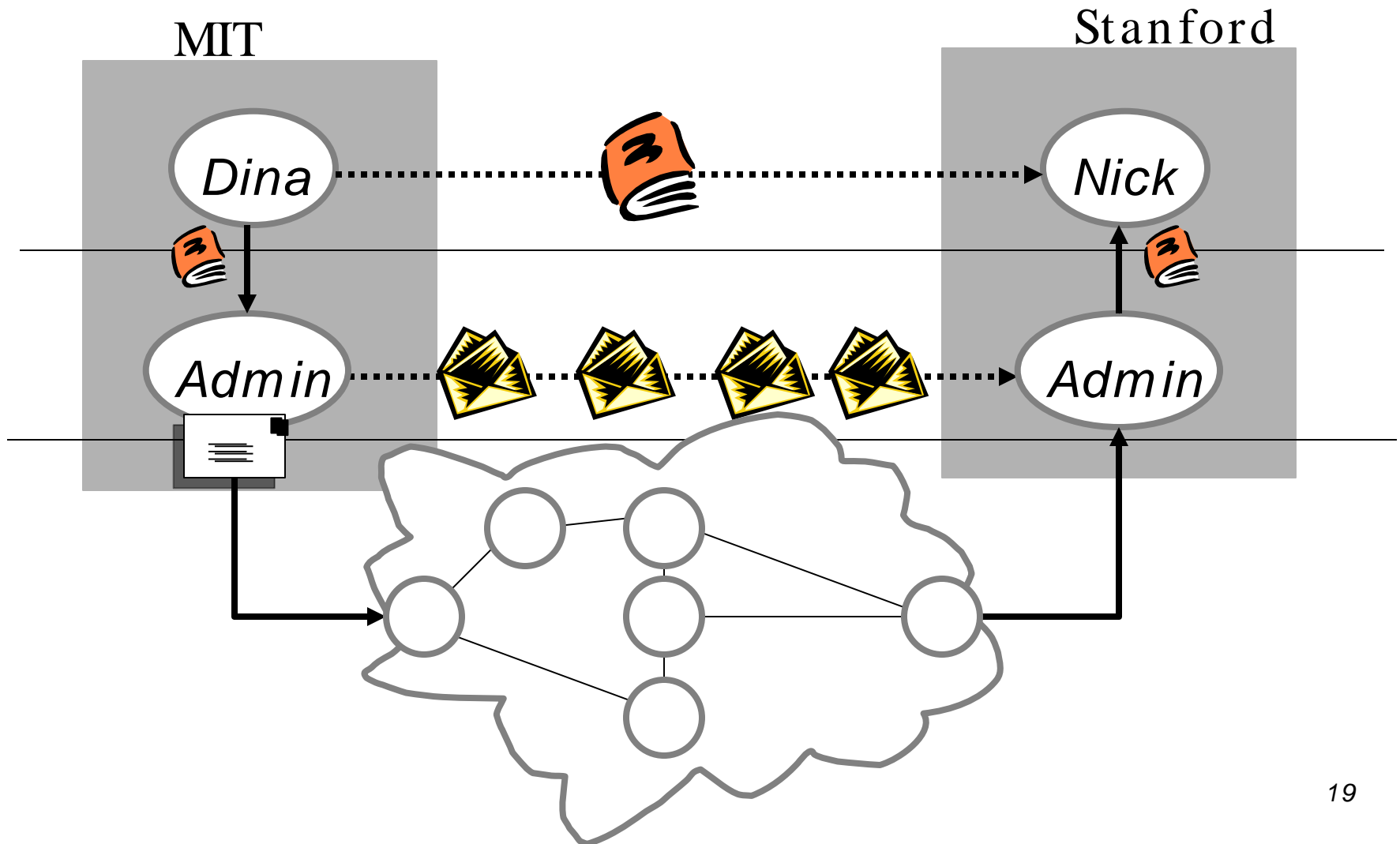


Best Effort Service

Analogy: the mail system

Internet's Best Effort Service

The mail system



Characteristics of the mail system

Each envelope is individually routed

No time guarantee for delivery

No guarantee of delivery in sequence

No guarantee of delivery at all!

Things get lost

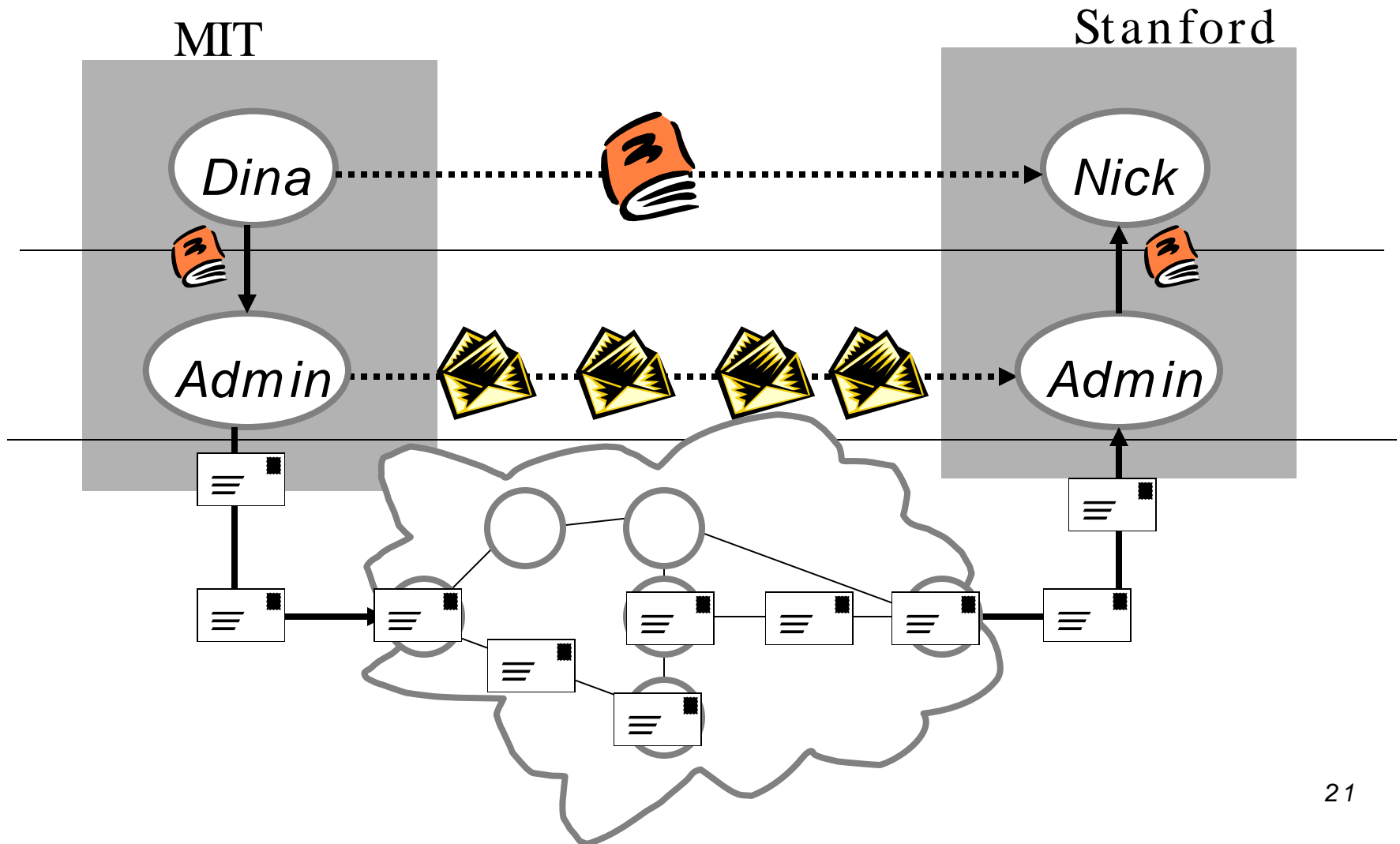
How can we acknowledge delivery?

Retransmission

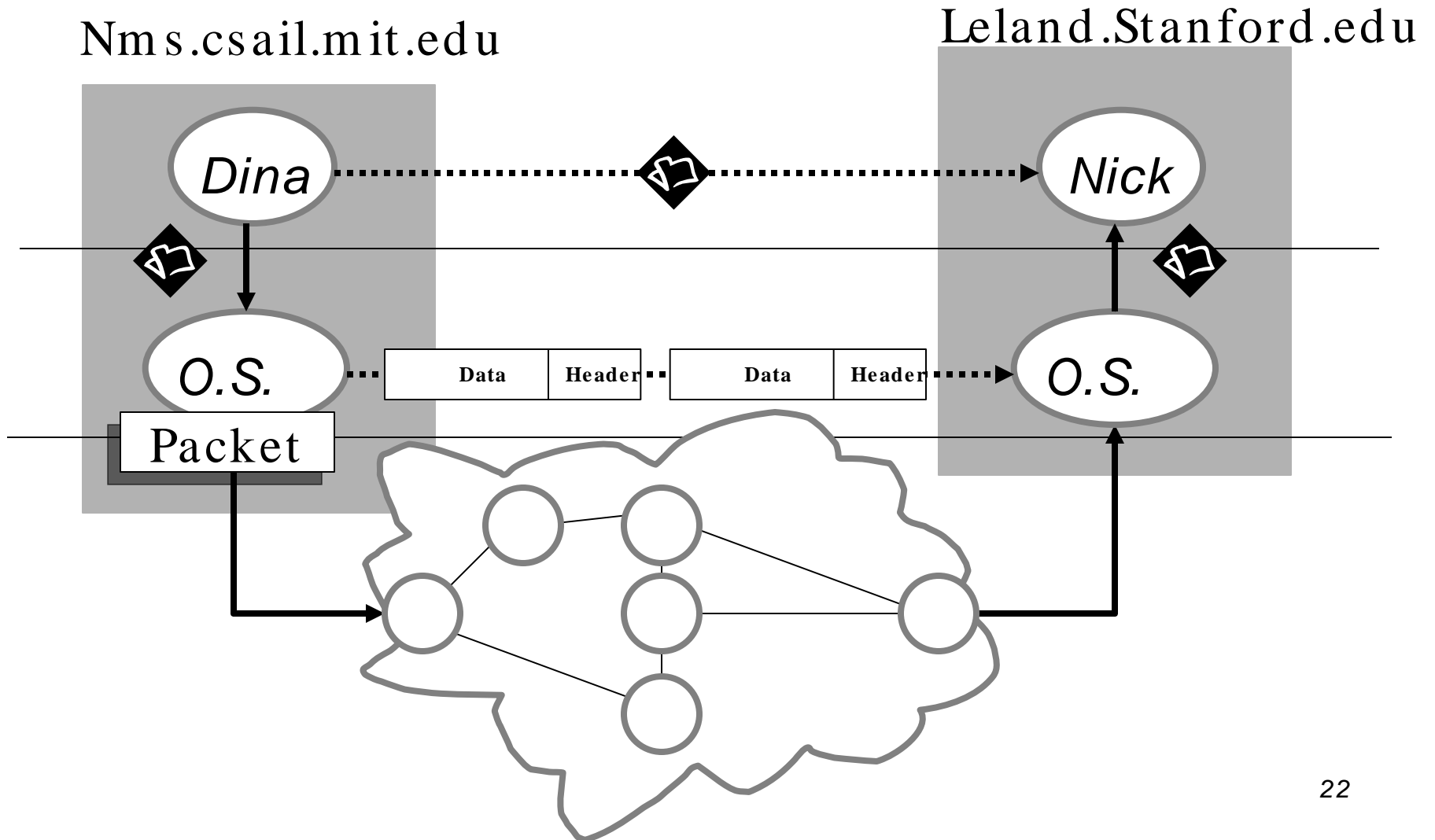
How to determine when to retransmit? Timeout?

*If message is re-sent too soon
duplicates*

The mail system



The Internet



Characteristics of the Internet

Each packet is individually routed

No time guarantee for delivery

No guarantee of delivery in sequence

No guarantee of delivery at all!

Things get lost

Acknowledgements

Retransmission

How to determine when to retransmit? Timeout?

*If packet is re-transmitted too soon
duplicate*

Best Effort

No Guarantees:

Variable Delay (jitter)

Variable rate

Packet loss

Duplicates

Reordering

(notes also state maximum packet length)

Differences Between Circuit & Packet Switching

<i>Circuit-switching</i>	<i>Packet-Switching</i>
<i>Guaranteed capacity</i>	<i>No guarantees (best effort)</i>
<i>Capacity is wasted if data is bursty</i>	<i>More efficient</i>
<i>Before sending data establishes a path</i>	<i>Send data immediately</i>
<i>All data in a single flow follow one path</i>	<i>Different packets might follow different paths</i>
<i>No reordering; constant delay; no pkt drops</i>	<i>Packets may be reordered, delayed, or dropped</i>

This Lecture

We learned how to share the network infrastructure between many connections/flows

We also learned about the implications of the sharing scheme (circuit or packet switching) on the service that the traffic receives